

Submodular Functions, Optimization, and Applications to Machine Learning

— Spring Quarter, Lecture 2 —

http://www.ee.washington.edu/people/faculty/bilmes/classes/ee596b_spring_2016/

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$$f(A) + f(B) \geq f(A \cup B) + f(A \cap B)$$

$$= f(A) + 2f(C) + f(B) - f(A) + f(C) + f(B) - f(A \cap B)$$



Cumulative Outstanding Reading

- Read chapter 1 from Fujishige's book.

Announcements, Assignments, and Reminders

- Weekly Office Hours: Mondays, 3:30-4:30, or by skype or google hangout (set up meeting via our our discussion board (https://canvas.uw.edu/courses/1039754/discussion_topics)).

Class Road Map - IT-I

- L1(3/28): Motivation, Applications, & Basic Definitions
- L2(3/30):
- L3(4/4):
- L4(4/6):
- L5(4/11):
- L6(4/13):
- L7(4/18):
- L8(4/20):
- L9(4/25):
- L10(4/27):
- L11(5/2):
- L12(5/4):
- L13(5/9):
- L14(5/11):
- L15(5/16):
- L16(5/18):
- L17(5/23):
- L18(5/25):
- L19(6/1):
- L20(6/6): Final Presentations maximization.

Finals Week: June 6th-10th, 2016.

Two Equivalent Submodular Definitions

Definition 2.2.1 (submodular concave)

A function $f : 2^V \rightarrow \mathbb{R}$ is **submodular** if for any $A, B \subseteq V$, we have that:

$$f(A) + f(B) \geq f(A \cup B) + f(A \cap B) \quad (2.8)$$

An alternate and (as we will soon see) equivalent definition is:

Definition 2.2.2 (diminishing returns)

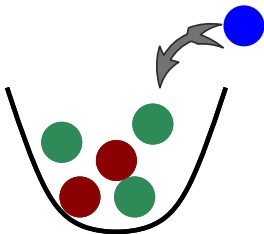
A function $f : 2^V \rightarrow \mathbb{R}$ is **submodular** if for any $A \subseteq B \subset V$, and $v \in V \setminus B$, we have that:

$$f(A \cup \{v\}) - f(A) \geq f(B \cup \{v\}) - f(B) \quad (2.9)$$

The incremental “value”, “gain”, or “cost” of v decreases (diminishes) as the context in which v is considered grows from A to B .

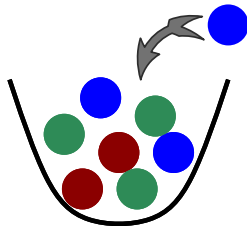
Example Submodular: Number of Colors of Balls in Urns

- Consider an urn containing colored balls. Given a set S of balls, $f(S)$ counts the number of distinct colors in S .



Initial value: 2 (colors in urn).

New value with added blue ball: 3



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- Submodularity: Incremental Value of Object Diminishes in a Larger Context (diminishing returns).
- Thus, f is submodular.

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Definition 2.2.2 (supermodular (improving returns))

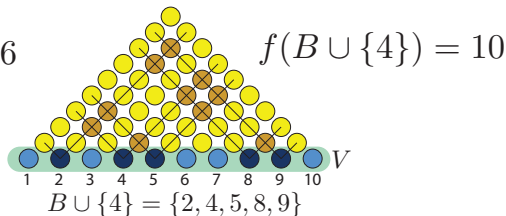
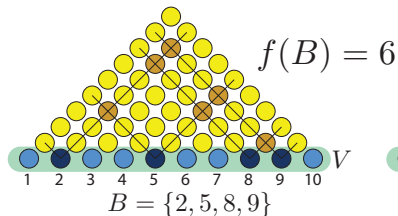
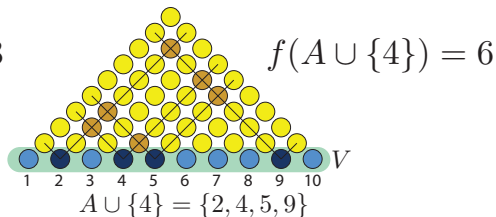
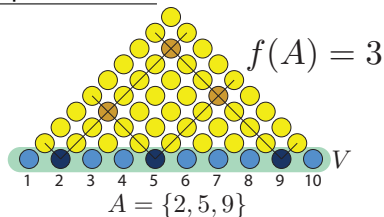
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- Incremental “value”, “gain”, or “cost” of v increases (improves) as the context in which v is considered grows from A to B .
- A function f is submodular iff $-f$ is supermodular.
- If f both submodular and supermodular, then f is said to be **modular**, and $f(A) = c + \sum_{a \in A} f(a)$ (often $c = 0$).

Example Supermodular: Number of Balls with Two Lines

Given ball pyramid, bottom row V is size $n = |V|$. For subset $S \subseteq V$ of bottom-row balls, draw 45° and 135° diagonal lines from each $s \in S$. Let $f(S)$ be number of non-bottom-row balls with two lines $\Rightarrow f(S)$ is supermodular.



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- Set cover, supply and demand side economies of scale,

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 - Also, we can “relax” a problem to a submodular one where it can be efficiently solved and offer a bounded quality solution.
 - Non-submodular problems can be analyzed via submodularity.

Many different functions are submodular!

- We will see many applications of submodularity in machine learning.
- On next set of slides, we will state (without proof, for now) that many of the functions are submodular (or supermodular).
- In subsequent lectures, we will start showing how to prove submodularity.

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 Answer: submodular maximization.

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Answer: submodular maximization.
- How do we choose the smallest set S that maintains a given degree of diversity? Constrained minimization (i.e., $\min |A|$ s.t. $f(A) \geq \alpha$).
- Random sample has probability of poorly representing normally underrepresented groups.

Extractive Document Summarization

- The figure below represents the sentences of a document



Extractive Document Summarization

- We extract sentences (green) as a summary of the full document

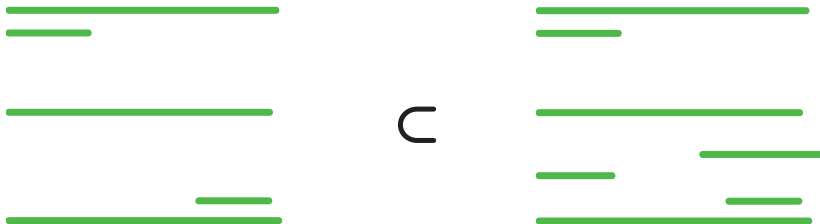


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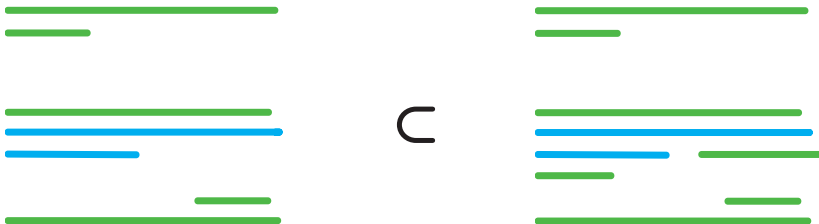
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- diminishing returns** \leftrightarrow **submodularity**

Large image collections need to be summarized

Many images, also that have a higher level gestalt than just a few, want a summary (subset of images) to represent the diversity in the large image set.

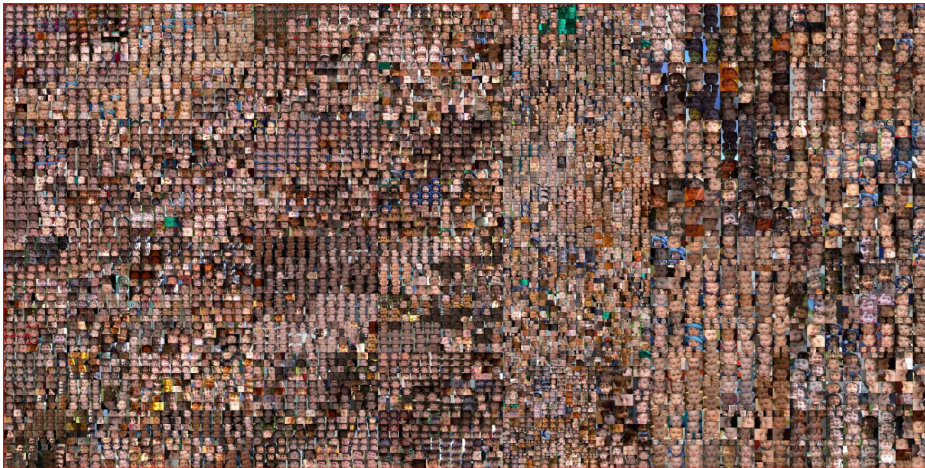


Image Summarization

10×10 image collection:



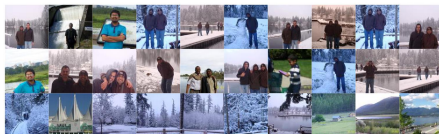
3 good summaries (diverse):



3 ok summaries:



3 poor summaries (redundant):



Variable Selection in Classification/Regression

- Let Y be a random variable we wish to accurately predict based on at most $n = |V|$ observed measurement variables $(X_1, X_2, \dots, X_n) = X_V$ in a probability model $\Pr(Y, X_1, X_2, \dots, X_n)$.

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- The mutual information function $f(A) = I(Y; X_A)$ is defined as:

$$I(Y; X_A) = \sum_{y, x_A} \Pr(y, x_A) \log \frac{\Pr(y, x_A)}{\Pr(y) \Pr(x_A)} = H(Y) - H(Y|X_A) \quad (2.1)$$

$$= H(X_A) - H(X_A|Y) = H(X_A) + H(Y) - H(X_A, Y) \quad (2.2)$$

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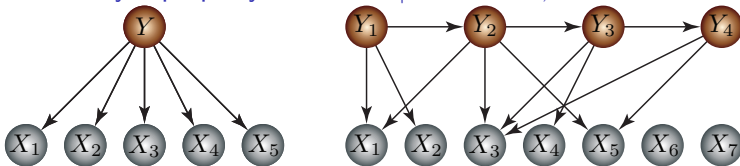
$$= H(X_A) - H(X_A|Y) = H(X_A) + H(Y) - H(X_A, Y) \quad (2.2)$$

- Applicable in pattern recognition, also in sensor coverage problem, where Y is whatever question we wish to ask about environment.

Information Gain and Feature Selection

in Pattern Classification: Naïve Bayes

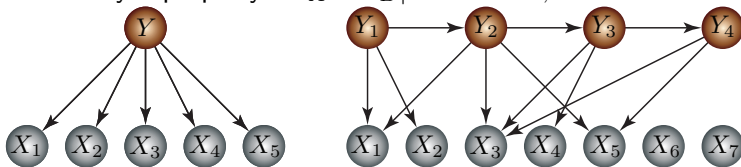
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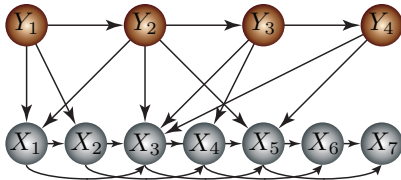
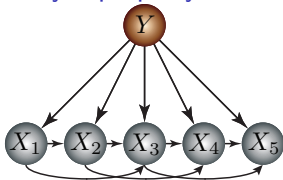
- When $X_A \perp\!\!\!\perp X_B | Y$ for all A, B (the Naïve Bayes assumption holds), then

$$f(A) = I(Y; X_A) = H(X_A) - H(X_A | Y) = H(X_A) - \sum_{a \in A} H(X_a | Y) \quad (2.3)$$

is submodular (submodular minus modular).

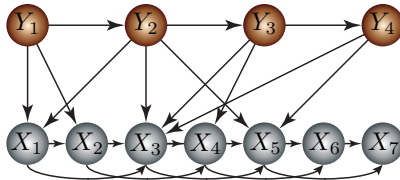
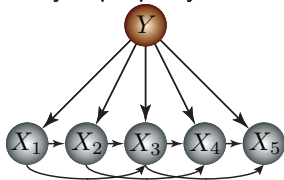
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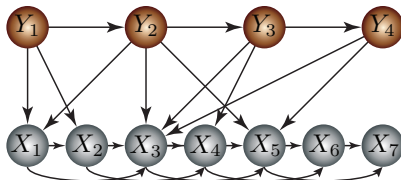
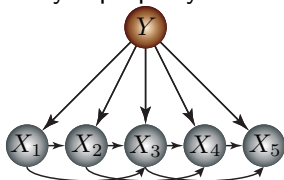
- $f(A)$ naturally expressed as a difference of two submodular functions

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- Alternatively, when Naïve Bayes assumption is false, we can make a submodular approximation (Peng-2005). E.g., functions of the form:

$$f(A) = \sum_{a \in A} I(X_a; Y) - \lambda \sum_{a, a' \in A} I(X_a; X_{a'} | Y) \quad (2.5)$$

where $\lambda \geq 0$ is a tradeoff constant.

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- $R_{Z,A}^2$'s minimizing parameters, for a given A , can be easily computed ($R_{Z,A}^2 = b_A^\top (C_A^{-1})^\top b_A$ when $\text{Var}Z = 1$, where $b_i = \text{Cov}(Z, X_i)$ and $C = E[(X - E[X])^\top (X - E[X])]$ is the covariance matrix).

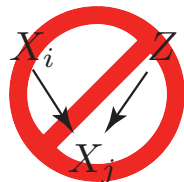
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$$R_{Z,A}^2 = \frac{\text{Var}(Z) - E[(Z - \tilde{Z}_A)^2]}{\text{Var}(Z)} \quad (2.6)$$

- $R_{Z,A}^2$'s minimizing parameters, for a given A , can be easily computed ($R_{Z,A}^2 = b_A^\top (C_A^{-1})^\top b_A$ when $\text{Var}Z = 1$, where $b_i = \text{Cov}(Z, X_i)$ and $C = E[(X - E[X])^\top (X - E[X])]$ is the covariance matrix).
- When there are no “suppressor” variables (essentially, no v-structures that converge on X_j with parents X_i and Z), then

$$f(A) = R_{Z,A}^2 = b_A^\top (C_A^{-1})^\top b_A \quad (2.7)$$



is a submodular function (so the greedy algorithm gives the $1 - 1/e$ guarantee). (Das&Kempe).

Data Subset Selection

- Suppose we are given a large data set $\mathcal{D} = \{x_i\}_{i=1}^n$ of n data items $V = \{v_1, v_2, \dots, v_n\}$ and we wish to choose a subset $A \subset V$ of items that is good in some way (e.g., a summary).

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- Example: U might be a set of textual features (e.g., ngrams), and $m_u(v)$ is the number of ngrams of type u in sentence v . E.g., if a document consists of the sentence

$v = \text{“Whenever I go to New York City, I visit the New York City museum.”}$

then $m_{\text{the}}(v) = 1$ while $m_{\text{New York City}}(v) = 2$.

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- $f(X)$ measures X 's ability to represent set of features U as measured by $m_u(X)$, with diminishing returns function g , and importance weights α_u .

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- Let $p = \{p_u\}_{u \in U}$ be a desired probability distribution over features (i.e., $\sum_u p_u = 1$ and $p_u \geq 0$ for all $u \in U$).

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- Consider the KL-divergence between these two distributions:

$$D(p || \{\bar{m}_u(X)\}_{u \in U}) = \sum_{u \in U} p_u \log p_u - \sum_{u \in U} p_u \log(\bar{m}_u(X)) \quad (2.11)$$

$$\begin{aligned} &= \sum_{u \in U} p_u \log p_u - \sum_{u \in U} p_u \log(m_u(X)) + \log(m(X)) \\ &= -H(p) + \log m(X) - \sum_{u \in U} p_u \log(m_u(X)) \end{aligned} \quad (2.12)$$

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- The objective once again, treating entropy $H(p)$ as a constant,

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- Hence the KL-divergence, seen as a function of X , i.e., $f(X) = D(p||\{\bar{m}_u(X)\})$ is quite naturally represented as a **difference of submodular functions**.
- Alternatively, if we define (Shinohara, 2014)

$$g(X) \triangleq \log m(X) - D(p||\{\bar{m}_u(X)\}) = \sum_{u \in U} p_u \log(m_u(X)) \quad (2.14)$$

we have a **submodular function** g that represents a combination of its quantity of X via its features (i.e., $\log m(X)$) and its feature distribution closeness to some distribution p (i.e., $D(p||\{\bar{m}_u(X)\})$).

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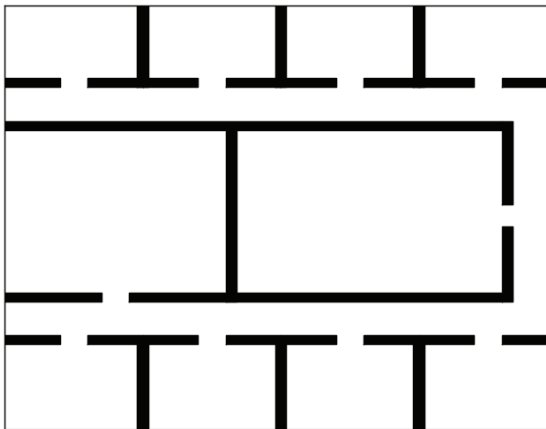
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- Environment could be a floor of a building, water network, monitored ecological preservation.

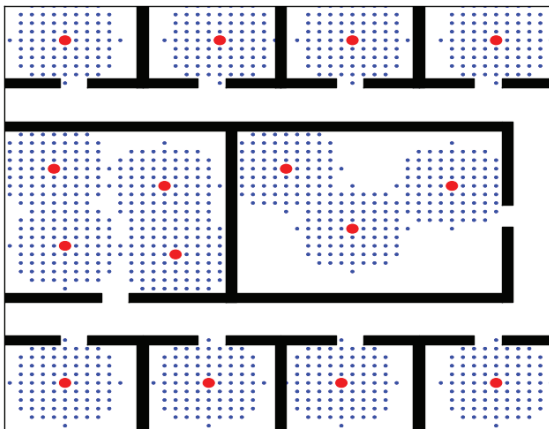
Sensor Placement within Buildings

- An example of a room layout. Should be possible to determine temperature at all points in the room. Sensors cannot sense beyond wall (thick black line) boundaries.



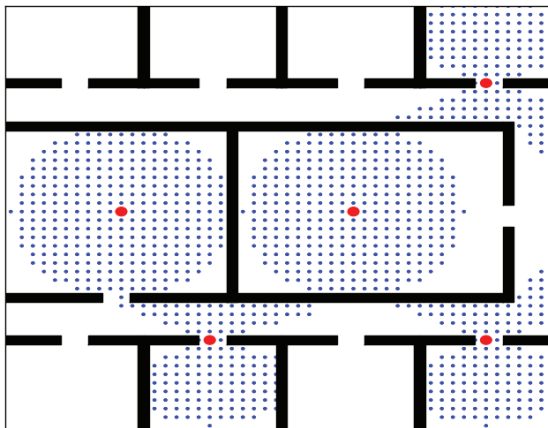
Sensor Placement within Buildings

- Example sensor placement using small range cheap sensors (located at red dots).



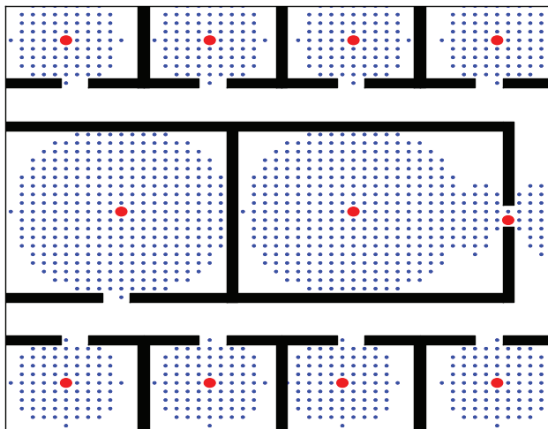
Sensor Placement within Buildings

- Example sensor placement using longer range expensive sensors (located at red dots).



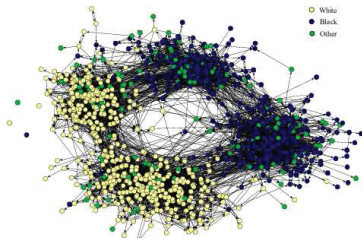
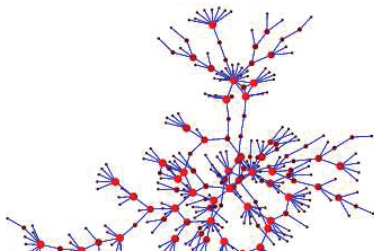
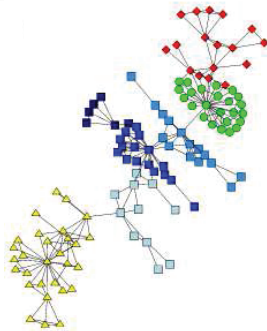
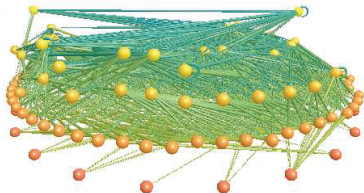
Sensor Placement within Buildings

- Example sensor placement using mixed range sensors (located at red dots).



Social Networks

(from Newman, 2004). Clockwise from top left: 1) predator-prey interactions, 2) scientific collaborations, 3) sexual contact, 4) school friendships.



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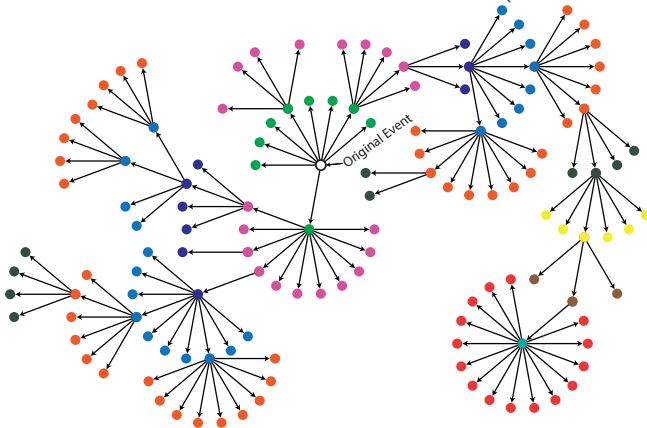
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- Which is a better model?

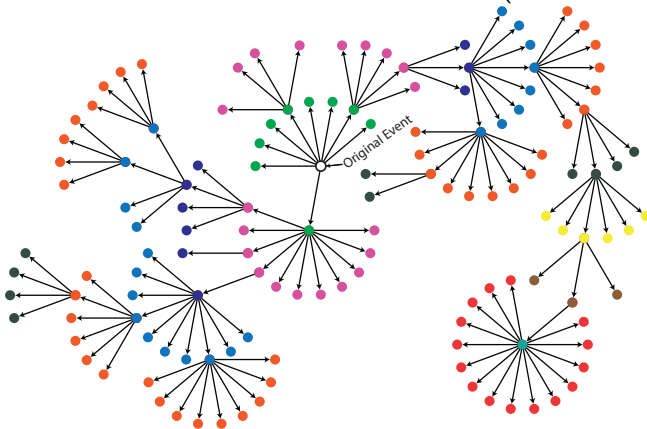
Information Cascades, Diffusion Networks

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- Goal: How to find the most influential sources, the ones that often set off cascades, which are like large “waves” of information flow?

Diffusion Networks

Where are they useful?

- **Information propagation:** when blogs or news stories break, and creates an information cascade over multiple other blogs/newspapers/magazines.
- **Viral marketing:** What is the pattern of trendsetters that cause an individual to purchase a product?
- **Epidemiology:** who gets sick from whom? What is the infection network of such links? Given finite supply of vaccine, who to inoculate to protect overall population (cut the network)?
 - Infer the connectivity of a network (memes, purchase decisions, viruses, etc.) based only on diffusion traces (the time that each node is "infected")?
 - How to find the most likely tree or graph?

A model of influence in social networks

- Given a graph $G = (V, E)$, each $v \in V$ corresponds to a person, to each v we have an activation function $f_v : 2^V \rightarrow [0, 1]$ dependent only on its neighbors. I.e., $f_v(A) = f_v(A \cap \Gamma(v))$.
- Goal - Viral Marketing: find a small subset $S \subseteq V$ of individuals to directly influence, and thus indirectly influence the greatest number of possible other individuals (via the social network G).
- Define function $f : 2^V \rightarrow \mathbb{Z}^+$ to model the ultimate influence of an initial infected nodes S . Use following iterative process; at each step:
 - Given previous set of infected nodes S that have not yet had their chance to infect their neighbors,
 - activate new nodes $v \in V \setminus S$ if $f_v(S \cap \Gamma_v) \geq U[0, 1]$, where $U[0, 1]$ is a uniform random number between 0 and 1, and Γ_v are the neighbors of v .
- For many f_v (including simple linear functions, and where f_v is submodular itself), we can show f is submodular (Kempe, Kleinberg, Tardos 1993).

Graphical Model Structure Learning

- A probability distribution on binary vectors $p : \{0, 1\}^V \rightarrow [0, 1]$:

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- This can be viewed as a discrete optimization problem on the potential (undirected) **edges** of the graph $V \times V$.

Graphical Models: Learning Tree Distributions

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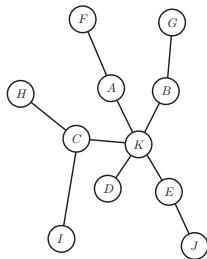
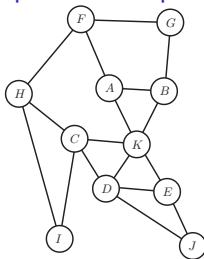
minimize
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subject to

$$D(p||p_t)$$

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$T = (V, F)$ is a tree



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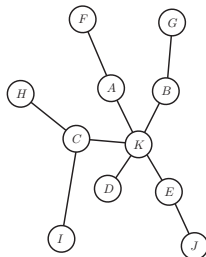
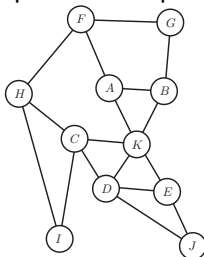
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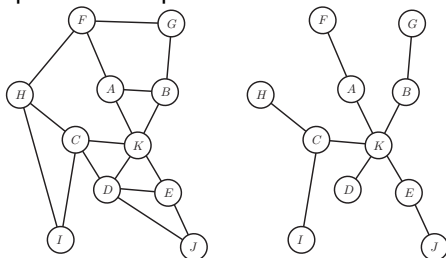
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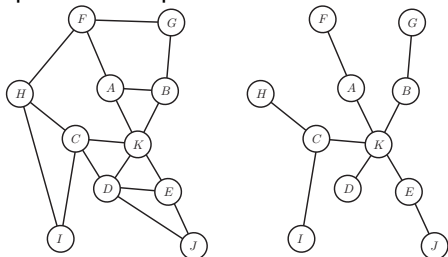
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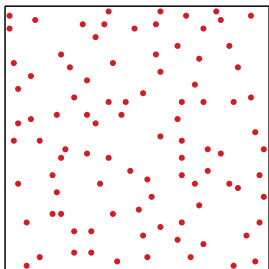
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- Then finding the maximum weight base of the matroid is solved by the greedy algorithm, and also finds the optimal tree (Chow & Liu, 1968)

Determinantal Point Processes (DPPs)

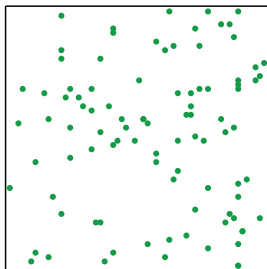
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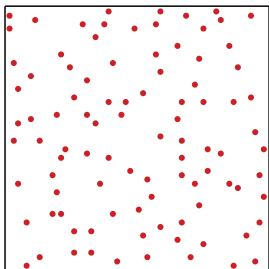


Independent

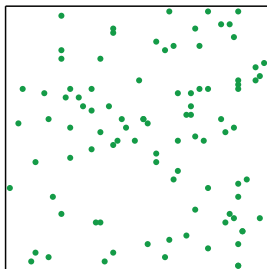
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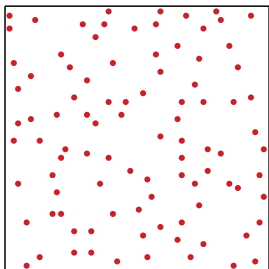
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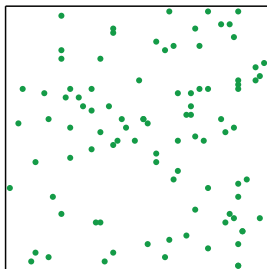
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- More “diverse” or “complex” samples are given higher probability.

DPPs and log-submodular probability distributions

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- Therefore, a DPP is a log-submodular probability distribution.

Graphical Models and fast MAP Inference

- Given distribution that factors w.r.t. a graph:

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- Many approximate inference strategies utilize additional factorization assumptions (e.g., mean-field, variational inference, expectation propagation, etc).
- Can we do exact MAP inference in polynomial time regardless of the tree-width, without even knowing the tree-width?

Order-two (edge) graphical models

- Given G let $p \in \mathcal{F}(G, \mathcal{M}^{(f)})$ such that we can write the **global energy** $E(x)$ as a sum of **unary** and **pairwise** potentials:

$$E(x) = \sum_{v \in V(G)} e_v(x_v) + \sum_{(i,j) \in E(G)} e_{ij}(x_i, x_j) \quad (2.21)$$

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- Further, say that $D_{X_v} = \{0, 1\}$ (binary), so we have binary random vectors distributed according to $p(x)$.
- Thus, $x \in \{0, 1\}^V$, and finding MPE solution is setting some of the variables to 0 and some to 1, i.e.,

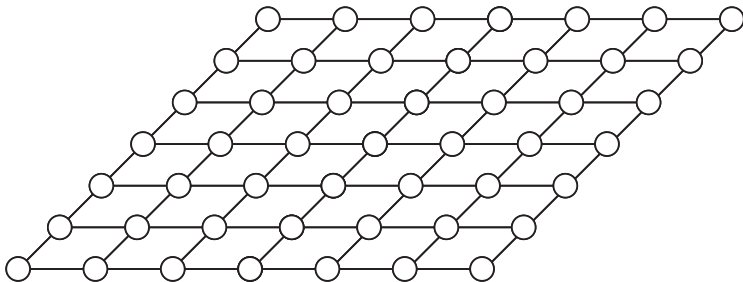
$$\min_{x \in \{0,1\}^V} E(x) \quad (2.22)$$

MRF example

Markov random field

$$\log p(x) \propto \sum_{v \in V(G)} e_v(x_v) + \sum_{(i,j) \in E(G)} e_{ij}(x_i, x_j) \quad (2.23)$$

When G is a 2D grid graph, we have



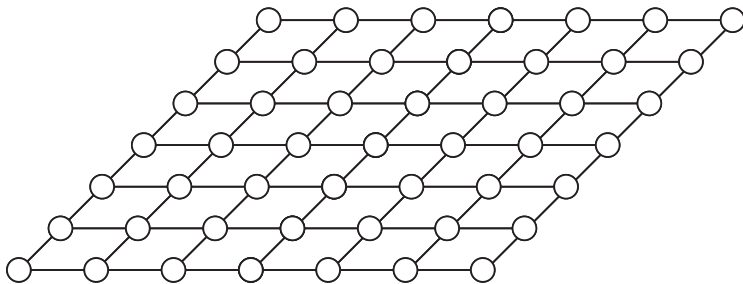
Create an auxiliary graph

- We can create auxiliary graph G_a that involves two new “terminal” nodes s and t and all of the original “non-terminal” nodes $v \in V(G)$.
- The non-terminal nodes represent the original random variables $x_v, v \in V$.
- Starting with the original grid-graph amongst the vertices $v \in V$, we connect each of s and t to all of the original nodes.
- I.e., we form $G_a = (V \cup \{s, t\}, E + \cup_{v \in V} ((s, v) \cup (v, t)))$.

Transformation from graphical model to auxiliary graph

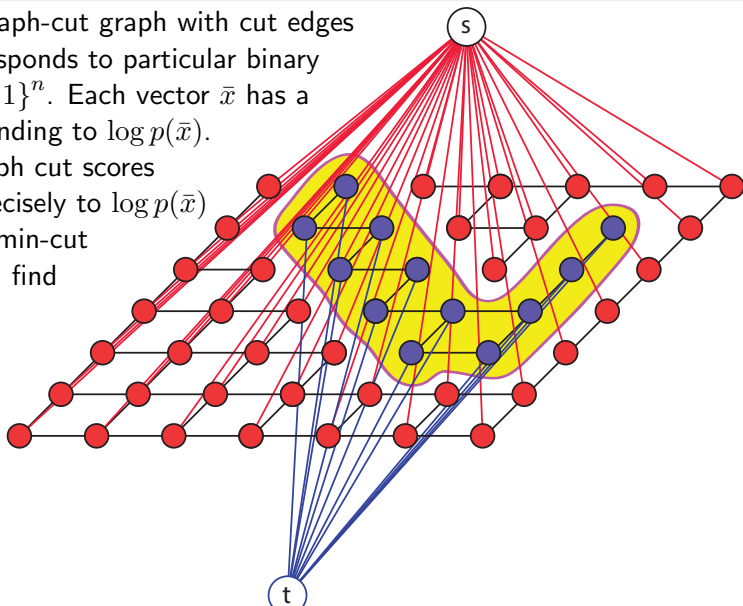
Original 2D-grid graphical model G and energy function

$E(x) = \sum_{v \in V(G)} e_v(x_v) + \sum_{(i,j) \in E(G)} e_{ij}(x_i, x_j)$ needing to be minimized over $x \in \{0, 1\}^V$. Recall, tree-width is $O(\sqrt{|V|})$.



Transformation from graphical model to auxiliary graph

Augmented graph-cut graph with cut edges removed corresponds to particular binary vector $\bar{x} \in \{0, 1\}^n$. Each vector \bar{x} has a score corresponding to $\log p(\bar{x})$. When can graph cut scores correspond precisely to $\log p(\bar{x})$ in a way that min-cut algorithms can find minimum of energy $E(x)$?



Setting of the weights in the auxiliary cut graph

- Any graph cut corresponds to a vector $\bar{x} \in \{0, 1\}^n$.
- If weights of all edges, except those involving terminals s and t , are non-negative, graph cut computable in polynomial time via max-flow (many algorithms, e.g., Edmonds&Karp $O(nm^2)$ or $O(n^2m \log(nC))$; Goldberg&Tarjan $O(nm \log(n^2/m))$, see Schrijver, page 161).
- If weights are set correctly in the cut graph, and if edge functions e_{ij} satisfy certain properties, then graph-cut score corresponding to \bar{x} can be made equivalent to $E(x) = \log p(\bar{x}) + \text{const.}$.
- Hence, poly time graph cut, can find the optimal MPE assignment, regardless of the graphical model's tree-width!
- In general, finding MPE is an NP-hard optimization problem.

Submodular potentials

submodularity is what allows graph cut to find exact solution

- Edge functions must be **submodular** (in the binary case, equivalent to “associative”, “attractive”, “regular”, “Potts”, or “ferromagnetic”):
for all $(i, j) \in E(G)$, must have:

$$e_{ij}(0, 1) + e_{ij}(1, 0) \geq e_{ij}(1, 1) + e_{ij}(0, 0) \quad (2.31)$$

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- As a set function, this is the same as:

$$f(X) = \sum_{\{i,j\} \in \mathcal{E}(G)} f_{i,j}(X \cap \{i, j\}) \quad (2.32)$$

which is submodular if each of the $f_{i,j}$'s are submodular!

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- A special case of more general submodular functions – unconstrained submodular function minimization is solvable in polytime.

On log-supermodular vs. log-submodular distributions

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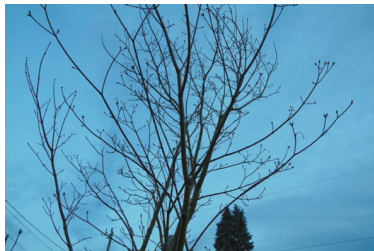
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- Log-submodular distributions:

$$\log \Pr(x) = f(x) + \text{const.} \quad (2.34)$$

where f is submodular. MAP or high-probable assignments should be “diverse”, or “complex”, or “covering”, like in determinantal point processes.

Shrinking bias in graph cut image segmentation



What does graph-cut based image segmentation do with elongated structures (top) or contrast gradients (bottom)?

Shrinking bias in graph cut image segmentation



Addressing shrinking bias with edge submodularity

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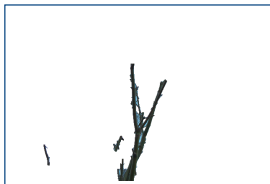
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- \Rightarrow cooperative-cut (Jegelka & B., 2011).

Graph-cut vs. cooperative-cut comparisons

Graph Cut

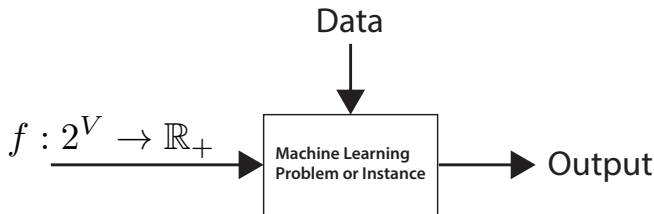
Cooperative Cut



(Jegelka&Bilmes,'11). There are fast algorithms for solving as well.

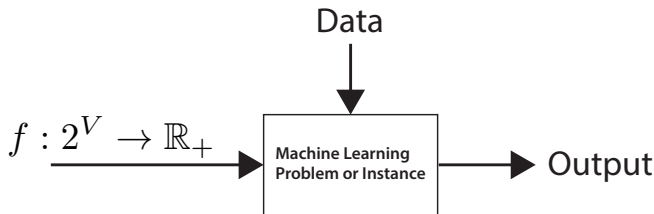
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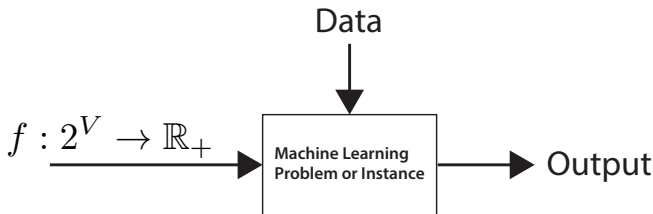
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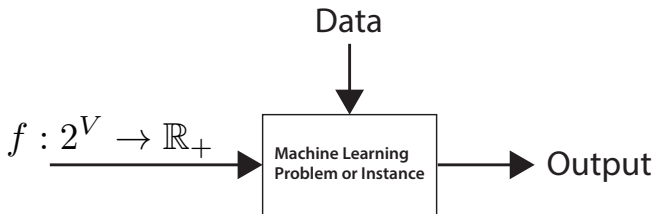
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- \mathbb{S} is a submodular cone since submodularity is closed under non-negative (conic) combinations.
- 2^n -dimensional since for certain $f \in \mathbb{S}$, there exists $f_\epsilon \in \mathbb{R}^{2^n}$ having no zero elements with $f + f_\epsilon \in \mathbb{S}$ (more on problem sets).

Supervised Machine Learning

From F. Bach

- We are given n samples of observed data $(x_i, y_i) \in \mathbb{R}^p \times \mathbb{R}$, $i \in [n]$.
 - Response vector $y = (y_1, \dots, y_n)^\top \in \mathbb{R}^n$
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$$\min_{w \in \mathbb{R}^p} \frac{1}{n} \sum_{i=1}^n \ell(y_i, w^\top x_i) + \lambda \Omega(w) = \min_{w \in \mathbb{R}^p} L(y, Xw) + \lambda \Omega(w) \quad (2.37)$$

where $\ell(\cdot)$ is a loss function (e.g., squared error) and $\Omega(w)$ is a (perhaps sparse) norm.

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- When data has multiple (k) responses, $y = (y^1, \dots, y^k) \in \mathbb{R}^{n \times k}$, we get:

$$\min_{w^1, \dots, w^k \in \mathbb{R}^n} \sum_{j=1}^k \{L(y^j, Xw^j) + \lambda \Omega(w^j)\} \quad (2.38)$$

Dictionary Learning and Selection

- When only the multiple responses $y = (y^1, \dots, y^k) \in \mathbb{R}^{n \times k}$ are observed, we get either **dictionary learning**

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- This is a subset selection problem, and the regularizer $\Omega(\cdot)$ is critical (could be structured sparse convex norm, via Lovász extension!).

Norms, sparse norms, and computer vision

- Common norms include p -norm $\Omega(w) = \|w\|_p = (\sum_{i=1}^p w_i^p)^{1/p}$
- 1-norm promotes sparsity (prefer solutions with zero entries).
- Image denoising, **total variation** is useful, norm takes form:

$$\Omega(w) = \sum_{i=2}^N |w_i - w_{i-1}| \quad (2.41)$$

related to Lovász extension of a graph-cut submodular function.

- Points of difference should be “sparse” (frequently zero).



(Rodriguez,
2009)

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- Ex: total variation is the Lovász-extension of graph cut

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- and two notions of “information amongst a collection of sets”:

$$I_f(S_1; S_2; \dots; S_k) = \sum_{i=1}^k f(S_i) - f(S_1 \cup S_2 \cup \dots \cup S_k) \quad (2.46)$$

$$I'_f(S_1; S_2; \dots; S_k) = \sum_{A \subseteq \{1, 2, \dots, k\}} (-1)^{|A|+1} f\left(\bigcup_{j \in A} S_j\right) \quad (2.47)$$

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- Hence, family of clustering algorithms parameterized by f .

Is Submodular Maximization Just Clustering?

- 1 Clustering objectives often NP-hard and inapproximable, submodular maximization is approximable for any submodular function.
- 2 To have guarantee, clustering typically needs metricity, submodularity parameterized via any non-negative pairwise values.
- 3 Clustering often requires separate process to choose representatives within each cluster. Submodular max does this automatically. Can also do submodular data partitioning (like clustering).
- 4 Submodular max covers clustering objectives such as k -medoids.
- 5 Can learn submodular functions (hence, learn clustering objective).
- 6 We can choose quality guarantee for any submodular function via submodular set cover (only possible for some clustering algorithms).
- 7 Submodular max with constraints, ensures representatives are feasible (e.g., knapsack, matroid independence, combinatorial, submodular level set, etc.)
- 8 Submodular functions may be more general than clustering objectives (submodularity allows high-order interactions between elements).

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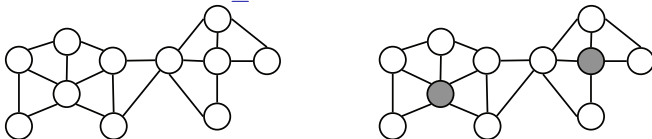
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- Adaptive active learning: choose a policy whereby we choose an $i_1 \in V$, get the label y_{i_1} , choose another $i_2 \in V$, get label y_{i_2} , where each chose can be based on previously acquired labels.

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- Often, getting y is time-consuming, expensive, and error prone (manual transcription, Amazon Turk, etc.)
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- Adaptive active learning: choose a policy whereby we choose an $i_1 \in V$, get the label y_{i_1} , choose another $i_2 \in V$, get label y_{i_2} , where each chose can be based on previously acquired labels.
- Semi-supervised (transductive) learning: Once we have $\{y_i\}_{i \in S}$, infer the remaining labels $\{y_i\}_{i \in V \setminus S}$.

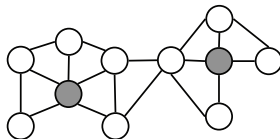
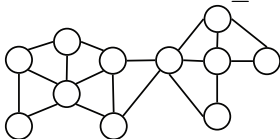
Active Transductive Semi-Supervised Learning

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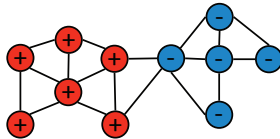
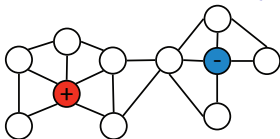


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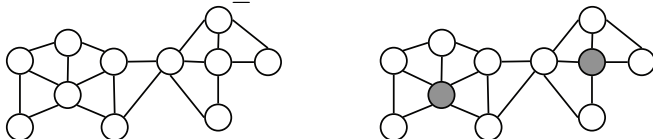


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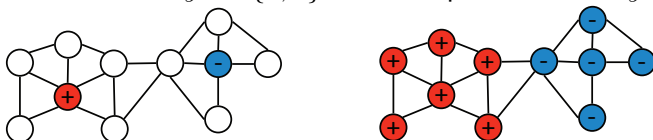


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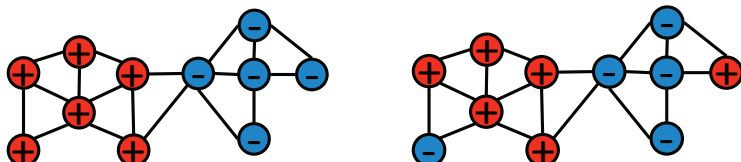
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- Learner suffers loss $\|\hat{y} - y\|_1$, where y is truth. Below, $\|\hat{y} - y\|_1 = 2$.



Choosing labels: how to select L

- Consider the following objective

$$\Psi(L) = \min_{T \subseteq V \setminus L: T \neq \emptyset} \frac{\Gamma(T)}{|T|} \quad (2.48)$$

where $\Gamma(T) = I_f(T; V \setminus T) = f(T) + f(V \setminus T) - f(V)$ is an arbitrary symmetric submodular function (e.g., graph cut value between T and $V \setminus T$, or combinatorial mutual information).

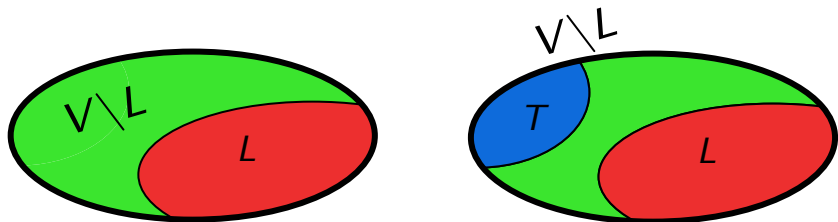
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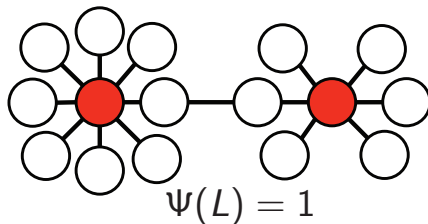
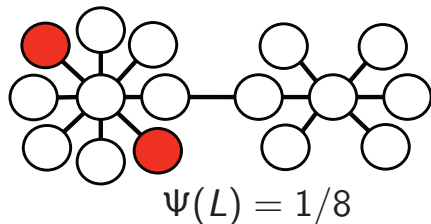
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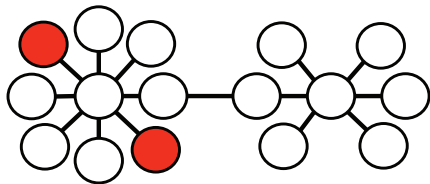
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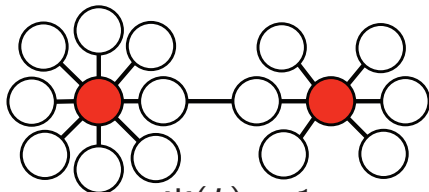
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- This suggests choosing (bounded cost) L that maximizes $\Psi(L)$.

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- In graph cut case, this is standard min-cut (Blum & Chawla 2001) approach to semi-supervised learning.

Generalized Error Bound

Theorem 2.6.1 (Guillory & B., '11)

For any symmetric submodular $\Gamma(S)$, assume \hat{y} minimizes $\Gamma(Y(\hat{y}))$ subject to $\hat{y}_L = y_L$. Then

$$\|\hat{y} - y\|_1 \leq 2 \frac{\Gamma(Y(y))}{\Psi(L)} \quad (2.50)$$

where $y \in \{0, 1\}^V$ are the true labels.

- All is defined in terms of the symmetric submodular function Γ (need not be graph cut), where:

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- $\Gamma(T) = I_f(T; V \setminus T) = f(S) + f(V \setminus S) - f(V)$ determined by arbitrary submodular function f , different error bound for each.
- Joint algorithm is “parameterized” by a submodular function f .

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- General: Hamming, Recall, Precision, Cond. MI, Sq. Hamming, etc.

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- *Balcan & Harvey (2011)*: submodular function learning problem from a learning theory perspective, given a distribution on subsets. Negative result is that can't approximate in this setting to within a constant factor.
- But can we learn a subclass, perhaps non-negative weighted mixtures of submodular components?

Structured Learning of Submodular Mixtures

- Constraints specified in inference form:

$$\underset{\mathbf{w}, \xi_t}{\text{minimize}} \quad \frac{1}{T} \sum_t \xi_t + \frac{\lambda}{2} \|\mathbf{w}\|^2 \quad (2.54)$$

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- If loss is supermodular, this is a difference-of-submodular (DS) function optimization.

Structured Prediction: Subgradient Learning

- Solvable with simple sub-gradient descent algorithm using structured variant of hinge-loss (Taskar, 2004).
- Loss-augmented inference is either submodular optimization (Lin & B. 2012) or DS optimization (Tschitschek, Iyer, & B. 2014).

Algorithm 1: Subgradient descent learning

Input : $S = \{(\mathbf{x}^{(t)}, \mathbf{y}^{(t)})\}_{t=1}^T$ and a learning rate sequence $\{\eta_t\}_{t=1}^T$.

1 $w_0 = 0$;

2 **for** $t = 1, \dots, T$ **do**

3 Loss augmented inference: $\mathbf{y}_t^* \in \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}_t} \mathbf{w}_{t-1}^\top \mathbf{f}_t(\mathbf{y}) + \ell_t(\mathbf{y})$;

4 Compute the subgradient: $\mathbf{g}_t = \lambda \mathbf{w}_{t-1} + \mathbf{f}_t(\mathbf{y}^*) - \mathbf{f}_t(\mathbf{y}^{(t)})$;

5 Update the weights: $\mathbf{w}_t = \mathbf{w}_{t-1} - \eta_t \mathbf{g}_t$;

Return : the averaged parameters $\frac{1}{T} \sum_t \mathbf{w}_t$.

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- Hence, rather than minimize $E(x)$ (hard), we can minimize $E_f(x) \geq E(x)$ (relatively easy), which is an upper bound.

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- For some variable selection problems, can get bounds of the form:

$$\text{Solution} \geq \left(1 - \frac{1}{e^{\gamma_{U^*,k}}}\right) \text{OPT} \quad (2.59)$$

where U^* is the solution set of a variable selection algorithm.

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- f is submodular if $\gamma_{U,k} \geq 1$ for all U and k .
- For some variable selection problems, can get bounds of the form:

$$\text{Solution} \geq \left(1 - \frac{1}{e^{\gamma_{U^*,k}}}\right) \text{OPT} \quad (2.59)$$

where U^* is the solution set of a variable selection algorithm.

- This gradually get worse as we move away from an objective being submodular (see Das & Kempe, 2011).

Submodular Analysis for Non-Submodular Problems

- Sometimes the quality of solutions to non-submodular problems can be analyzed via submodularity.
- For example, “deviation from submodularity” can be measured using the **submodularity ratio** (Das & Kempe):

$$\gamma_{U,k}(f) = \min_{L \subseteq U, S: |S| \leq k, S \cap L = \emptyset} \frac{\sum_{s \in S} f(x|L)}{f(S|L)} \quad (2.58)$$

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- This gradually get worse as we move away from an objective being submodular (see Das & Kempe, 2011).
- Other analogous concepts: **curvature** of a submodular function, and also the **submodular degree**.

Monge Matrices

- $m \times n$ matrices $C = [c_{ij}]_{ij}$ are called Monge matrices if they satisfy the Monge property, namely:

$$c_{ij} + c_{rs} \leq c_{is} + c_{rj} \quad (2.60)$$

for all $1 \leq i < r \leq m$ and $1 \leq j < s \leq n$.

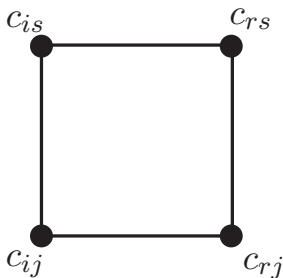
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- Consider four elements of the matrix:



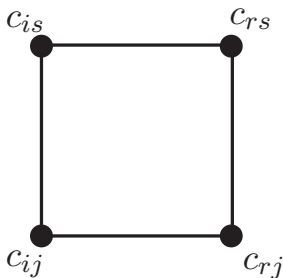
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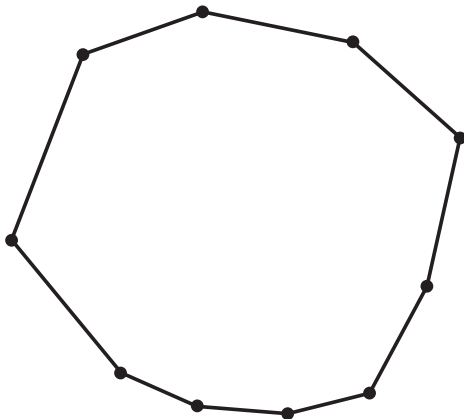
- Consider four elements of the matrix:



- Useful for speeding up certain dynamic programming problems.

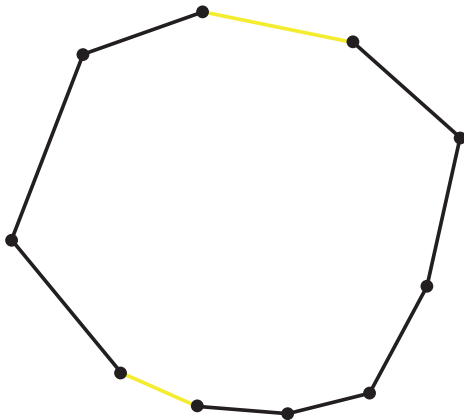
Monge Matrices

- Can generate a Monge matrix from a convex polygon - delete two segments, then separately number vertices on each chain. Distances c_{ij} satisfy Monge property (or quadrangle inequality).



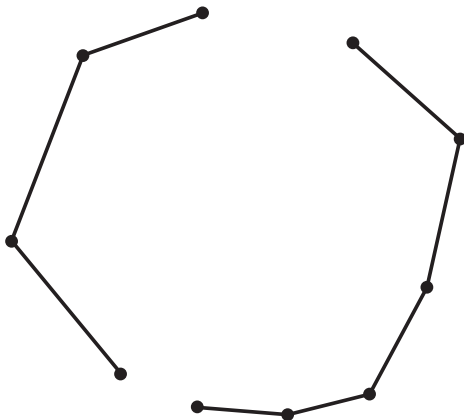
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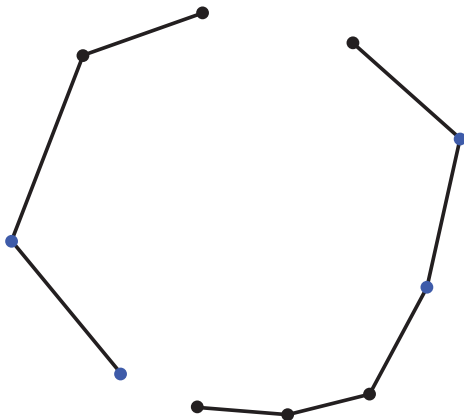
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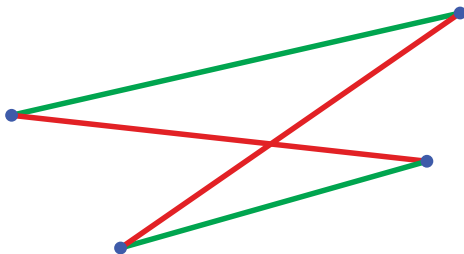
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Monge Matrices

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Example Submodular: Entropy from Information Theory

- Entropy is submodular. Let V be the index set of a set of random variables, then the function

$$f(A) = H(X_A) = - \sum_{x_A} p(x_A) \log p(x_A) \quad (2.61)$$

is submodular.

- Proof: conditioning reduces entropy. With $A \subseteq B$ and $v \notin B$,

$$H(X_v|X_B) = H(X_{B+v}) - H(X_B) \quad (2.62)$$

$$\leq H(X_{A+v}) - H(X_A) = H(X_v|X_A) \quad (2.63)$$

Information Theory: Block Coding

- Given a set of random variables $\{X_i\}_{i \in V}$ indexed by set V , how do we partition them so that we can best block-code them within each block.

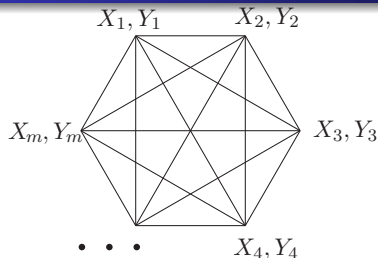
Information Theory: Block Coding

- Given a set of random variables $\{X_i\}_{i \in V}$ indexed by set V , how do we partition them so that we can best block-code them within each block.
- I.e., how do we form $S \subseteq V$ such that $I(X_S; X_{V \setminus S})$ is as small as possible, where $I(X_A; X_B)$ is the mutual information between random variables X_A and X_B , i.e.,

$$I(X_A; X_B) = H(X_A) + H(X_B) - H(X_A, X_B) \quad (2.64)$$

and $H(X_A) = -\sum_{x_A} p(x_A) \log p(x_A)$ is the joint entropy of the set X_A of random variables.

Information Theory: Network Communication



- A network of senders/receivers
- Each sender X_i is trying to communicate simultaneously with each receiver Y_i (i.e., for all i , X_i is sending to $\{Y_i\}_i$)
- The X_i are **not** necessarily independent.
- Communication rates from i to j are $R^{(i \rightarrow j)}$ to send message $W^{(i \rightarrow j)} \in \{1, 2, \dots, 2^{nR^{(i \rightarrow j)}}\}$.
- Goal: necessary and sufficient conditions for achievability.
- I.e., can we find functions f such that any rates must satisfy

$$\forall S \subseteq V, \quad \sum_{i \in S, j \in V \setminus S} R^{(i \rightarrow j)} \leq f(S) \quad (2.65)$$

- Special cases MAC (Multi-Access Channel) for communication over $p(y|x_1, x_2)$ and Slepian-Wolf compression (independent compression of X and Y but at joint rate $H(X, Y)$).

Example Submodular: Entropy from Information Theory

- Alternate Proof: Conditional mutual Information is always non-negative.
- Given $A, B \subseteq V$, consider conditional mutual information quantity:

$$\begin{aligned} I(X_{A \setminus B}; X_{B \setminus A} | X_{A \cap B}) &= \sum_{x_{A \cup B}} p(x_{A \cup B}) \log \frac{p(x_{A \setminus B}, x_{B \setminus A} | x_{A \cap B})}{p(x_{A \setminus B} | x_{A \cap B}) p(x_{B \setminus A} | x_{A \cap B})} \\ &= \sum_{x_{A \cup B}} p(x_{A \cup B}) \log \frac{p(x_{A \cup B}) p(x_{A \cap B})}{p(x_A) p(x_B)} \geq 0 \quad (2.66) \end{aligned}$$

then

$$\begin{aligned} I(X_{A \setminus B}; X_{B \setminus A} | X_{A \cap B}) \\ = H(X_A) + H(X_B) - H(X_{A \cup B}) - H(X_{A \cap B}) \geq 0 \quad (2.67) \end{aligned}$$

so entropy satisfies

$$H(X_A) + H(X_B) \geq H(X_{A \cup B}) + H(X_{A \cap B}) \quad (2.68)$$

Example Submodular: Mutual Information

- Also, symmetric mutual information is submodular,

$$f(A) = I(X_A; X_{V \setminus A}) = H(X_A) + H(X_{V \setminus A}) - H(X_V) \quad (2.69)$$

Note that $f(A) = H(X_A)$ and $\bar{f}(A) = H(X_{V \setminus A})$, and adding submodular functions preserves submodularity (which we will see quite soon).

Two Equivalent Submodular Definitions

Definition 2.11.1 (submodular concave)

A function $f : 2^V \rightarrow \mathbb{R}$ is **submodular** if for any $A, B \subseteq V$, we have that:

$$f(A) + f(B) \geq f(A \cup B) + f(A \cap B) \quad (2.8)$$

An alternate and (as we will soon see) equivalent definition is:

Definition 2.11.2 (diminishing returns)

A function $f : 2^V \rightarrow \mathbb{R}$ is **submodular** if for any $A \subseteq B \subset V$, and $v \in V \setminus B$, we have that:

$$f(A \cup \{v\}) - f(A) \geq f(B \cup \{v\}) - f(B) \quad (2.9)$$

The incremental “value”, “gain”, or “cost” of v decreases (diminishes) as the context in which v is considered grows from A to B .

Subadditive Definitions

Definition 2.11.1 (subadditive)

A function $f : 2^V \rightarrow \mathbb{R}$ is subadditive if for any $A, B \subseteq V$, we have that:

$$f(A) + f(B) \geq f(A \cup B) \quad (2.70)$$

This means that the “whole” is less than the sum of the parts.

Two Equivalent Supermodular Definitions

Definition 2.11.1 (supermodular)

A function $f : 2^V \rightarrow \mathbb{R}$ is **supermodular** if for any $A, B \subseteq V$, we have that:

$$f(A) + f(B) \leq f(A \cup B) + f(A \cap B) \quad (2.8)$$

Definition 2.11.2 (supermodular (improving returns))

A function $f : 2^V \rightarrow \mathbb{R}$ is **supermodular** if for any $A \subseteq B \subset V$, and $v \in V \setminus B$, we have that:

$$f(A \cup \{v\}) - f(A) \leq f(B \cup \{v\}) - f(B) \quad (2.9)$$

- Incremental “value”, “gain”, or “cost” of v increases (improves) as the context in which v is considered grows from A to B .
- A function f is submodular iff $-f$ is supermodular.
- If f both submodular and supermodular, then f is said to be **modular**, and $f(A) = c + \sum_{a \in A} f(a)$ (often $c = 0$).

Superadditive Definitions

Definition 2.11.2 (superadditive)

A function $f : 2^V \rightarrow \mathbb{R}$ is superadditive if for any $A, B \subseteq V$, we have that:

$$f(A) + f(B) \leq f(A \cup B) \quad (2.71)$$

- This means that the “whole” is greater than the sum of the parts.

Superadditive Definitions

Definition 2.11.2 (superadditive)

A function $f : 2^V \rightarrow \mathbb{R}$ is superadditive if for any $A, B \subseteq V$, we have that:

$$f(A) + f(B) \leq f(A \cup B) \quad (2.71)$$

- This means that the “whole” is greater than the sum of the parts.
- In general, submodular and subadditive (and supermodular and superadditive) are different properties.

Modular Definitions

Definition 2.11.3 (modular)

A function that is both submodular and supermodular is called **modular**

If f is a modular function, then for any $A, B \subseteq V$, we have

$$f(A) + f(B) = f(A \cap B) + f(A \cup B) \quad (2.72)$$

In modular functions, elements do not interact (or cooperate, or compete, or influence each other), and have value based only on singleton values.

Proposition 2.11.4

If f is modular, it may be written as

$$f(A) = f(\emptyset) + \sum_{a \in A} \left(f(\{a\}) - f(\emptyset) \right) \quad (2.73)$$

Modular Definitions

Proof.

We inductively construct the value for $A = \{a_1, a_2, \dots, a_k\}$.

For $k = 2$,

$$f(a_1) + f(a_2) = f(a_1, a_2) + f(\emptyset) \quad (2.74)$$

$$\text{implies } f(a_1, a_2) = f(a_1) - f(\emptyset) + f(a_2) - f(\emptyset) + f(\emptyset) \quad (2.75)$$

then for $k = 3$,

$$f(a_1, a_2) + f(a_3) = f(a_1, a_2, a_3) + f(\emptyset) \quad (2.76)$$

$$\text{implies } f(a_1, a_2, a_3) = f(a_1, a_2) - f(\emptyset) + f(a_3) - f(\emptyset) + f(\emptyset) \quad (2.77)$$

$$= f(\emptyset) + \sum_{i=1}^3 (f(a_i) - f(\emptyset)) \quad (2.78)$$

and so on ...



Complement function

Given a function $f : 2^V \rightarrow \mathbb{R}$, we can find a complement function $\bar{f} : 2^V \rightarrow \mathbb{R}$ as $\bar{f}(A) = f(V \setminus A)$ for any A .

Proposition 2.11.5

\bar{f} is submodular if f is submodular.

Proof.

$$\bar{f}(A) + \bar{f}(B) \geq \bar{f}(A \cup B) + \bar{f}(A \cap B) \quad (2.79)$$

follows from

$$f(V \setminus A) + f(V \setminus B) \geq f(V \setminus (A \cup B)) + f(V \setminus (A \cap B)) \quad (2.80)$$

which is true because $V \setminus (A \cup B) = (V \setminus A) \cap (V \setminus B)$ and $V \setminus (A \cap B) = (V \setminus A) \cup (V \setminus B)$. □