Submodular Functions, Optimization, and Applications to Machine Learning — Spring Quarter, Lecture 16 — <u>http://j.ee.washington.edu/~bilmes/classes/ee596b\_spring\_2014/</u>

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May 21st, 2014



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EE596b/Spring 2014/Submodularity - Lecture 16 - May 21st, 2014

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## Cumulative Outstanding Reading

- Good references for today: Schrijver-2003, Oxley-1992/2011, Welsh-1973, Goemans-2010, Cunningham-1984, Edmonds-1969, Choquet-1955, Grabisch/Marichal/Mesiar/Pap "Aggregation Functions", Lovász-1983, Bach-2011.
- Read Tom McCormick's overview paper on SFM http://people. commerce.ubc.ca/faculty/mccormick/sfmchap8a.pdf
- Read chapters 1 4 from Fujishige book.
- Matroid properties http: //www-math.mit.edu/~goemans/18433S09/matroid-notes.pdf
- Read lecture 14 slides on lattice theory at our web page (http://j. ee.washington.edu/~bilmes/classes/ee596b\_spring\_2014/)

Logistics

Review



• Weekly Office Hours: Wednesdays, 5:00-5:50, or by skype or google hangout (email me).

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Review

### Class Road Map - IT-I

- L1 (3/31): Motivation, Applications, & Basic Definitions
- L2: (4/2): Applications, Basic Definitions, Properties
- L3: More examples and properties (e.g., closure properties), and examples, spanning trees
- L4: proofs of equivalent definitions, independence, start matroids
- L5: matroids, basic definitions and examples
- L6: More on matroids, System of Distinct Reps, Transversals, Transversal Matroid, Matroid and representation
- L7: Dual Matroids, other matroid properties, Combinatorial Geometries
- L8: Combinatorial Geometries, matroids and greedy, Polyhedra, Matroid Polytopes,
- L9: From Matroid Polytopes to Polymatroids.
- L10: Polymatroids and Submodularity

- L11: More properties of polymatroids, SFM special cases
- L12: polymatroid properties, extreme points polymatroids,
- L13: sat, dep, supp, exchange capacity, examples
- L14: Lattice theory: partially ordered sets; lattices; distributive, modular, submodular, and boolean lattices; ideals and join irreducibles.
- L15: Supp, Base polytope, polymatroids and entropic Venn diagrams, exchange capacity,
- L16: minimum norm point algorithm and the lattice of minimizers of a submodular function, Lovasz extension
- L17:
- L18:
- L19:
- L20:

#### Finals Week: June 9th-13th, 2014.

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Review

### Summary of supp, sat, and dep

• For 
$$x \in P_f$$
,  $\operatorname{supp}(x) = \{e : x(e) \neq 0\} \subseteq \operatorname{sat}(x)$ 

• For  $x \in P_f$ , sat(x) (span, closure) is the maximal saturated (x-tight) set w.r.t. x. I.e., sat $(x) = \{e : e \in E, \forall \alpha > 0, x + \alpha \mathbf{1}_e \notin P_f\}$ . That is,

$$\mathsf{cl}(x) \stackrel{\text{def}}{=} \operatorname{sat}(x) \triangleq \bigcup \left\{ A : A \in \mathcal{D}(x) \right\}$$
(16.29)

$$= \bigcup \left\{ A : A \subseteq E, x(A) = f(A) \right\}$$
(16.30)

$$= \{e : e \in E, \forall \alpha > 0, x + \alpha \mathbf{1}_e \notin P_f\}$$
(16.31)

• For  $e \in \operatorname{sat}(x)$ , we have  $\operatorname{dep}(x, e) \subseteq \operatorname{sat}(x)$  (fundamental circuit) is the minimal (common) saturated (x-tight) set w.r.t. x containing e. I.e.,  $\operatorname{dep}(x, e) = \begin{cases} \bigcap \{A : e \in A \subseteq E, x(A) = f(A)\} & \text{if } e \in \operatorname{sat}(x) \\ \emptyset & \text{else} \end{cases}$ 

$$= \left\{ e' : \exists \alpha > 0, \text{ s.t. } x + \alpha (\mathbf{1}_e - \mathbf{1}_{e'}) \in P_f \right\}$$
(16.32)

Logistics

# A polymatroid function's polyhedron is a polymatroid.

#### Theorem 16.2.1

Let f be a submodular function defined on subsets of E. For any  $x \in \mathbb{R}^E$  , we have:

$$rank(x) = \max\left(y(E) : y \le x, y \in \mathbf{P}_f\right) = \min\left(x(A) + f(E \setminus A) : A \subseteq E\right)$$
(16.5)

If we take x to be zero, we get:

Corollary 16.2.2

Let f be a submodular function defined on subsets of E.  $x \in \mathbb{R}^{E}$ , we have:

$$rank(0) = \max(y(E) : y \le 0, y \in P_f) = \min(f(A) : A \subseteq E)$$
 (16.6)

Review

Review

### Multiple Polytopes associated with f



$$P_f^+ = P_f \cap \left\{ x \in \mathbb{R}^E : x \ge 0 \right\}$$
(16.5)

$$P_f = \left\{ x \in \mathbb{R}^E : x(S) \le f(S), \forall S \subseteq E \right\}$$
(16.6)

$$B_f = P_f \cap \left\{ x \in \mathbb{R}^E : x(E) = f(E) \right\}$$
(16.7)

### Min-Norm Point: Definition

• Restating what we saw before, we have:

 $\max\{y(E)|y \in P_f, y \le 0\} = \min\{f(X)|X \subseteq V\}$ (16.12)

• Consider the optimization:

minimize
$$\|x\|_2^2$$
(16.13a)subject to $x \in B_f$ (16.13b)

where  $B_f$  is the base polytope of submodular f, and  $\|x\|_2^2 = \sum_{e \in E} x(e)^2$  is the squared 2-norm. Let  $x^*$  be the optimal solution.

- Note,  $x^*$  is the unique optimal solution since we have a strictly convex objective over a set of convex constraints.
- $x^*$  is called the minimum norm point of the base polytope.

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### Min-Norm Point: Examples



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Min-Norm Point and SFM

### Min-Norm Point and Submodular Function Minimization

 $\bullet\,$  Given optimal solution  $x^*$  to the above, consider the quantities

$$y^* = x^* \land 0 = (\min(x^*(e), 0) | e \in E)$$
(16.1)

$$A_{-} = \{e : x^{*}(e) < 0\}$$
(16.2)

$$A_0 = \{e : x^*(e) \le 0\}$$
(16.3)

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• Thus, we immediately have that:

$$A_{-} \subseteq A_{0} \tag{16.4}$$

and that

$$x^*(A_-) = x^*(A_0) = y^*(A_-) = y^*(A_0)$$
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• It turns out, these quantities will solve the submodular function minimization problem, as we now show.

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- It turns out, these quantities will solve the submodular function minimization problem, as we now show.
- The proof is nice since it uses the tools we've been recently developing.

#### Theorem 16.3.1

Let  $y^*$ ,  $A_-$ , and  $A_0$  be as given. Then  $y^*$  is a maximizer of the l.h.s. of Eqn. (15.12). Moreover,  $A_-$  is the unique minimal minimizer of f and  $A_0$  is the unique maximal minimizer of f.

#### Proof.

• First note, since  $x^* \in B_f$ , we have  $x^*(E) = f(E)$ , meaning  $\operatorname{sat}(x^*) = E$ . Thus, we can consider any  $e \in E$  within  $\operatorname{dep}(x^*, e)$ .

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- Consider any pair (e, e') with  $e' \in dep(x^*, e)$  and  $e \in A_-$ . Then  $x^*(e) < 0$ , and  $\exists \alpha > 0$  s.t.  $x^* + \alpha \mathbf{1}_e \alpha \mathbf{1}_{e'} \in P_f$ .

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- Consider any pair (e, e') with  $e' \in dep(x^*, e)$  and  $e \in A_-$ . Then  $x^*(e) < 0$ , and  $\exists \alpha > 0$  s.t.  $x^* + \alpha \mathbf{1}_e \alpha \mathbf{1}_{e'} \in P_f$ .
- We have  $x^*(E) = f(E)$  and  $x^*$  is minimum in I2 sense. We have  $(x^* + \alpha \mathbf{1}_e \alpha \mathbf{1}_{e'}) \in P_f$ , and in fact

$$(x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'})(E) = x^*(E) + \alpha - \alpha = f(E)$$
(16.6)

so  $x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'} \in B_f$  also.

#### ... proof of Thm. 16.3.1 cont.

# • Then $(x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'})(E)$ = $x^*(E \setminus \{e, e'\}) + \underbrace{(x^*(e) + \alpha)}_{x^*_{new}(e)} + \underbrace{(x^*(e') - \alpha)}_{x^*_{new}(e')} = f(E).$

... proof of Thm. 16.3.1 cont.

• Then 
$$(x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'})(E)$$
  
=  $x^*(E \setminus \{e, e'\}) + \underbrace{(x^*(e) + \alpha)}_{x^*_{new}(e)} + \underbrace{(x^*(e') - \alpha)}_{x^*_{new}(e')} = f(E).$   
• Minimality of  $x^* \in B_f$  in [2 sense requires that, with such an  $\alpha > 0$ .

• Minimality of 
$$x^* \in D_f$$
 in 12 sense requires that, with such  $(x^*(e))^2 + (x^*(e'))^2 < (x^*_{\mathsf{new}}(e))^2 + (x^*_{\mathsf{new}}(e'))^2$ 

... proof of Thm. 16.3.1 cont.

• Then 
$$(x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'})(E)$$
  
=  $x^*(E \setminus \{e, e'\}) + \underbrace{(x^*(e) + \alpha)}_{x^*_{\text{new}}(e)} + \underbrace{(x^*(e') - \alpha)}_{x^*_{\text{new}}(e')} = f(E).$ 

• Minimality of  $x^* \in B_f$  in l2 sense requires that, with such an  $\alpha > 0$ ,  $\left(x^*(e)\right)^2 + \left(x^*(e')\right)^2 < \left(x^*_{\text{new}}(e)\right)^2 + \left(x^*_{\text{new}}(e')\right)^2$ 

• Given that  $e \in A_-$ ,  $x^*(e) < 0$ . Thus, if  $x^*(e') > 0$ , we could have  $(x^*(e) + \alpha)^2 + (x^*(e') - \alpha)^2 < (x^*(e))^2 + (x^*(e'))^2$ , contradicting the optimality of  $x^*$ .

... proof of Thm. 16.3.1 cont.

• Then 
$$(x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'})(E)$$
  
=  $x^*(E \setminus \{e, e'\}) + \underbrace{(x^*(e) + \alpha)}_{x^*_{\text{new}}(e)} + \underbrace{(x^*(e') - \alpha)}_{x^*_{\text{new}}(e')} = f(E).$ 

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- If  $x^*(e') = 0$ , we would have  $(x^*(e) + \alpha)^2 + (\alpha)^2 < (x^*(e))^2$ , for any  $0 < \alpha < |x^*(e)|$  (Exercise:), again contradicting the optimality of  $x^*$ .

... proof of Thm. 16.3.1 cont.

• Then 
$$(x^* + \alpha \mathbf{1}_e - \alpha \mathbf{1}_{e'})(E)$$
  
=  $x^*(E \setminus \{e, e'\}) + \underbrace{(x^*(e) + \alpha)}_{x^*_{\text{new}}(e)} + \underbrace{(x^*(e') - \alpha)}_{x^*_{\text{new}}(e')} = f(E).$ 

- Minimality of  $x^* \in B_f$  in l2 sense requires that, with such an  $\alpha > 0$ ,  $\left(x^*(e)\right)^2 + \left(x^*(e')\right)^2 < \left(x^*_{\mathsf{new}}(e)\right)^2 + \left(x^*_{\mathsf{new}}(e')\right)^2$
- Given that  $e \in A_-$ ,  $x^*(e) < 0$ . Thus, if  $x^*(e') > 0$ , we could have  $(x^*(e) + \alpha)^2 + (x^*(e') \alpha)^2 < (x^*(e))^2 + (x^*(e'))^2$ , contradicting the optimality of  $x^*$ .
- If  $x^*(e') = 0$ , we would have  $(x^*(e) + \alpha)^2 + (\alpha)^2 < (x^*(e))^2$ , for any  $0 < \alpha < |x^*(e)|$  (Exercise:), again contradicting the optimality of  $x^*$ .
- Thus, we must have  $x^*(e') < 0$  (strict negativity).

### Min-Norm Point and SFM

#### ... proof of Thm. 16.3.1 cont.

• Thus, for a pair (e, e') with  $e' \in dep(x^*, e)$  and  $e \in A_-$ , we have x(e') < 0 and hence  $e' \in A_-$ .

#### ... proof of Thm. 16.3.1 cont.

- Thus, for a pair (e, e') with  $e' \in dep(x^*, e)$  and  $e \in A_-$ , we have x(e') < 0 and hence  $e' \in A_-$ .
- Hence,  $\forall e \in A_{-}$ , we have  $dep(x^*, e) \subseteq A_{-}$ .

#### ... proof of Thm. 16.3.1 cont.

- Thus, for a pair (e, e') with  $e' \in dep(x^*, e)$  and  $e \in A_-$ , we have x(e') < 0 and hence  $e' \in A_-$ .
- Hence,  $\forall e \in A_-$ , we have  $dep(x^*, e) \subseteq A_-$ .
- A very similar argument can show that,  $\forall e \in A_0$ , we have  $dep(x^*, e) \subseteq A_0$ .

Lovász extension

### Min-Norm Point and SFM

#### ... proof of Thm. 16.3.1 cont.

• Therefore, we have  $\cup_{e \in A_-} \operatorname{dep}(x^*, e) = A_-$  and  $\cup_{e \in A_0} \operatorname{dep}(x^*, e) = A_0$ 

#### ... proof of Thm. 16.3.1 cont.

• Therefore, we have 
$$\cup_{e \in A_-} \operatorname{dep}(x^*, e) = A_-$$
 and  $\cup_{e \in A_0} \operatorname{dep}(x^*, e) = A_0$ 

• Ie.,  $\{\operatorname{dep}(x^*, e)\}_{e \in A_-}$  is cover for  $A_-$ , as is  $\{\operatorname{dep}(x^*, e)\}_{e \in A_0}$  for  $A_0$ .

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- Therefore, we have  $\cup_{e \in A_-} \operatorname{dep}(x^*, e) = A_-$  and  $\cup_{e \in A_0} \operatorname{dep}(x^*, e) = A_0$
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- $dep(x^*, e)$  is minimal tight set containing e, meaning  $x^*(dep(x^*, e)) = f(dep(x^*, e))$ , and since tight sets are closed under union, we have that  $A_-$  and  $A_0$  are also tight, meaning:

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$$x^*(A_-) = f(A_-) \tag{16.7}$$

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 $x^*(A_0) = f(A_0)$  (16.8)

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$$x^{*}(A_{-}) = x^{*}(A_{0}) = y^{*}(E) = y^{*}(A_{0}) + \underbrace{y^{*}(E \setminus A_{0})}_{=0}$$
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and therefore, all together we have

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$$x^*(A_-) = f(A_-) \tag{16.7}$$

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$$x^{*}(A_{-}) = x^{*}(A_{0}) = y^{*}(E) = y^{*}(A_{0}) + \underbrace{y^{*}(E \setminus A_{0})}_{=0}$$
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and therefore, all together we have

$$f(A_{-}) = f(A_{0}) = x^{*}(A_{-}) = x^{*}(A_{0}) = y^{*}(E)$$
 (16.10)

#### ... proof of Thm. 16.3.1 cont.

• Now,  $y^*$  is feasible for the l.h.s. of Eqn. (15.12).

#### ... proof of Thm. 16.3.1 cont.

Now, y\* is feasible for the l.h.s. of Eqn. (15.12). This follows since, we have y\* = x\* ∧ 0 ≤ 0, and since x\* ∈ B<sub>f</sub> ⊂ P<sub>f</sub>, and y\* ≤ x\* and P<sub>f</sub> is down-closed, we have that y\* ∈ P<sub>f</sub>.

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- Now,  $y^*$  is feasible for the l.h.s. of Eqn. (15.12). This follows since, we have  $y^* = x^* \land 0 \leq 0$ , and since  $x^* \in B_f \subset P_f$ , and  $y^* \leq x^*$  and  $P_f$  is down-closed, we have that  $y^* \in P_f$ .
- Also, for any  $y \in P_f$  with  $y \leq 0$  and for any  $X \subseteq E$ , we have  $y(E) \leq y(X) \leq f(X)$ .

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- Now, y\* is feasible for the l.h.s. of Eqn. (15.12). This follows since, we have y\* = x\* ∧ 0 ≤ 0, and since x\* ∈ B<sub>f</sub> ⊂ P<sub>f</sub>, and y\* ≤ x\* and P<sub>f</sub> is down-closed, we have that y\* ∈ P<sub>f</sub>.
- Also, for any  $y \in P_f$  with  $y \leq 0$  and for any  $X \subseteq E$ , we have  $y(E) \leq y(X) \leq f(X)$ .
- Hence, we have found a feasible for l.h.s. of Eqn. (15.12),  $y^* \leq 0$ ,  $y^* \in P_f$ , so  $y^*(E) \leq f(X)$  for all X.
### ... proof of Thm. 16.3.1 cont.

- Now,  $y^*$  is feasible for the l.h.s. of Eqn. (15.12). This follows since, we have  $y^* = x^* \land 0 \leq 0$ , and since  $x^* \in B_f \subset P_f$ , and  $y^* \leq x^*$  and  $P_f$  is down-closed, we have that  $y^* \in P_f$ .
- Also, for any  $y \in P_f$  with  $y \le 0$  and for any  $X \subseteq E$ , we have  $y(E) \le y(X) \le f(X)$ .
- Hence, we have found a feasible for l.h.s. of Eqn. (15.12),  $y^* \leq 0$ ,  $y^* \in P_f$ , so  $y^*(E) \leq f(X)$  for all X.
- So  $y^*(E) \le \min{\{f(X) | X \subseteq V\}}$ .

. . .

#### ... proof of Thm. 16.3.1 cont.

- Now, y\* is feasible for the l.h.s. of Eqn. (15.12). This follows since, we have y\* = x\* ∧ 0 ≤ 0, and since x\* ∈ B<sub>f</sub> ⊂ P<sub>f</sub>, and y\* ≤ x\* and P<sub>f</sub> is down-closed, we have that y\* ∈ P<sub>f</sub>.
- Also, for any  $y \in P_f$  with  $y \le 0$  and for any  $X \subseteq E$ , we have  $y(E) \le y(X) \le f(X)$ .
- Hence, we have found a feasible for l.h.s. of Eqn. (15.12),  $y^* \leq 0$ ,  $y^* \in P_f$ , so  $y^*(E) \leq f(X)$  for all X.
- So  $y^*(E) \le \min{\{f(X) | X \subseteq V\}}$ .
- Considering Eqn. (16.7), we have found sets  $A_{-}$  and  $A_{0}$  with tightness in Eqn. (15.12), meaning  $y^{*}(E) = f(A_{-}) = f(A_{0})$ .

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- Considering Eqn. (16.7), we have found sets  $A_{-}$  and  $A_{0}$  with tightness in Eqn. (15.12), meaning  $y^{*}(E) = f(A_{-}) = f(A_{0})$ .
- Hence,  $y^*$  is a maximizer of l.h.s. of Eqn. (15.12), and  $A_-$  and  $A_0$  are minimizers of f.

. . .

#### ... proof of Thm. 16.3.1 cont.

• Now, for any  $X \subset A_-$ , we have

$$f(X) \ge x^*(X) > x^*(A_-) = f(A_-)$$
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• And for any 
$$X \supset A_0$$
, we have

$$f(X) \ge x^*(X) > x^*(A_0) = f(A_0)$$
(16.12)

• Hence,  $A_{-}$  must be the unique minimal minimizer of f, and  $A_{0}$  is the unique maximal minimizer of f.

F16/38 (pg.42/127)

• So, if we have a procedure to compute the min-norm point computation, we can solve SFM.

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- An algorithm (by F. Wolfe) can find this min-norm point, essentially an active-set procedure for quadratic programming. It uses Edmonds's greedy algorithm to make it efficient.
- This is currently the best practical algorithm for general purpose submodular function minimization.
- But its underlying lower-bound complexity is unknown, although in practice its estimated empirical complexity runs anywhere from  $O(n^3)$  to  $O(n^{4.5})$  or so (see Jegelka, Lin, Bilmes (NIPS 2011)).

• Recall, that the set of minimizers of f forms a lattice.

F18/38 (pg.49/127)

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- In fact, with  $x^*$  the min-norm point, and  $A_-$  and  $A_0$  as defined above, we have the following theorem:

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- In fact, with  $x^*$  the min-norm point, and  $A_-$  and  $A_0$  as defined above, we have the following theorem:

Theorem 16.3.2

Let  $A \subseteq E$  be any minimizer of submodular f, and let  $x^*$  be the minimum-norm point. Then A has the form:

 $A = A_{-} \cup \bigcup_{a \in A_{m}} \operatorname{dep}(x^{*}, a)$ (16.13)

for some set  $A_m \subseteq A_0 \setminus A_{-}$ .

Min-Norm Point and SFM

### proof of Thm. 16.3.2.

• If A is a minimizer, then  $A_{-} \subseteq A \subseteq A_{0}$ , and  $f(A) = y^{*}(E)$  is the minimum valuation of f.

Min-Norm Point and SFM

- If A is a minimizer, then  $A_{-} \subseteq A \subseteq A_{0}$ , and  $f(A) = y^{*}(E)$  is the minimum valuation of f.
- But  $x^* \in P_f$ , so  $x^*(A) \leq f(A)$  and  $f(A) = x^*(A_-) \leq x^*(A)$  (or alternatively, just note that  $x^*(A_0 \setminus A) = 0$ ).

## Min-norm point and other minimizers of f

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- Hence,  $x^*(A) = x^*(A_-) = f(A)$  so that A is also a tight set for  $x^*$ .
- For any a ∈ A, A is a tight set containing a, and dep(x\*, a) is the minimal tight containing a.

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- For any  $a \in A$ , A is a tight set containing a, and  $dep(x^*, a)$  is the minimal tight containing a.
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- Hence,  $x^*(A) = x^*(A_-) = f(A)$  so that A is also a tight set for  $x^*$ .
- For any  $a \in A$ , A is a tight set containing a, and  $dep(x^*, a)$  is the minimal tight containing a.
- Hence, for any  $a \in A$ ,  $dep(x^*, a) \subseteq A$ .
- This means that  $\bigcup_{a\in A} \operatorname{dep}(x^*,a) = A.$
- Since  $A_{-} \subseteq A \subseteq A_{0}$ , then  $\exists A_{m} \subseteq A \setminus A_{-}$  such that

$$A = \bigcup_{a \in A_{-}} \operatorname{dep}(x^*, a) \cup \bigcup_{a \in A_{m}} \operatorname{dep}(x^*, a) = A_{-} \cup \bigcup_{a \in A_{m}} \operatorname{dep}(x^*, a)$$

## On a unique minimizer f

• Note that if f(e|A) > 0,  $\forall A \subseteq E$  and  $e \in E \setminus A$ , then we have  $A_{-} = A_0$  (there is one unique minimizer).

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# On a unique minimizer f

- Note that if f(e|A) > 0,  $\forall A \subseteq E$  and  $e \in E \setminus A$ , then we have  $A_{-} = A_{0}$  (there is one unique minimizer).
- On the other hand, if  $A_{-} = A_{0}$ , it does not imply f(e|A) > 0 for all  $A \subseteq E \setminus \{e\}$ .

# On a unique minimizer f

- Note that if f(e|A) > 0,  $\forall A \subseteq E$  and  $e \in E \setminus A$ , then we have  $A_{-} = A_{0}$  (there is one unique minimizer).
- On the other hand, if  $A_- = A_0$ , it does not imply f(e|A) > 0 for all  $A \subseteq E \setminus \{e\}$ .
- If  $A_- = A_0$  then certainly  $f(e|A_0) > 0$  for  $e \in E \setminus A_0$  and  $-f(e|A_0 \setminus \{e\}) > 0$  for all  $e \in A_0$ .



The next slide comes from lecture 12.

# Polymatroidal polyhedron and greedy

 Thus, restating the above results into a single complete theorem, we have a result very similar to what we saw for matroids (i.e., Theorem ??)

#### Theorem 16.4.1

Min-Norm Point and SFM

If  $f: 2^E \to \mathbb{R}_+$  is given, and P is a polytope in  $\mathbb{R}^E_+$  of the form  $P = \{x \in \mathbb{R}^E_+ : x(A) \le f(A), \forall A \subseteq E\}$ , then the greedy solution to the problem  $\max(wx: x \in P)$  is  $\forall w$  optimum iff f is monotone non-decreasing submodular (i.e., iff P is a polymatroid).

# Optimization over $P_f$

### • Consider the following optimization. Given $w \in \mathbb{R}^{E}$ ,

maximize	$w^{\intercal}x$	(16.14a)
subject to	$x \in P_f$	(16.14b)

# Optimization over $P_f$

Min-Norm Point and SFM

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maximize	$w^\intercal x$	(16.14a)
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• Since  $P_f$  is down closed, if  $\exists e \in E$  with w(e) < 0 then the solution above is unboundedly large.

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Min-Norm Point and SFM

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- The greedy algorithm will solve this, and the proof almost identical.

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  - $\begin{array}{ll} \text{maximize} & w^{\mathsf{T}}x & (16.14a)\\ \text{subject to} & x \in P_f & (16.14b) \end{array}$
- Since  $P_f$  is down closed, if  $\exists e \in E$  with w(e) < 0 then the solution above is unboundedly large. Hence, assume  $w \in \mathbb{R}^E_+$ .
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- Due to Theorem 15.5.2, any  $x \in P_f$  with  $x \notin B_f$  is dominated by  $x \leq y \in B_f$  which can only increase  $w^{\mathsf{T}}x \leq w^{\mathsf{T}}y$ .

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• Moreover, we can have  $w \in \mathbb{R}^E$  if we insist on  $x \in B_f$ .

## A continuous extension of f

• Consider again optimization problem. Given  $w \in \mathbb{R}^E$ ,

maximize	$w^\intercal x$	(16.16a)
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• We may consider this optimization problem a function  $\tilde{f}: \mathbb{R}^E \to \mathbb{R}$  of  $w \in \mathbb{R}^E$ , defined as:

 $\tilde{f}(w) = \max(wx : x \in P_f) \tag{16.17}$
Min-Norm Point and SFM

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$$\tilde{f}(w) = \max(wx : x \in P_f) \tag{16.17}$$

Hence, for any w, from the above theorem, we can compute the value of this function using the greedy algorithm (after of course checking for w ∈ ℝ<sup>E</sup><sub>+</sub>).

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#### A continuous extension of submodular f

• That is, given a submodular function f, a  $w \in \mathbb{R}^E$ , and defining  $E_i = \{e_1, e_2, \dots, e_i\}$  and where we choose the element order  $(e_1, e_2, \dots, e_m)$  based on decreasing w,so that  $w(e_1) \ge w(e_2) \ge \dots \ge w(e_m)$ , we have  $\tilde{f}(w)$ 

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Lovász extension

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- That is, given a submodular function f, a  $w \in \mathbb{R}^{E}$ , and defining  $E_i = \{e_1, e_2, \dots, e_i\}$  and where we choose the element order  $(e_1, e_2, \ldots, e_m)$  based on decreasing w, so that  $w(e_1) > w(e_2) > \cdots > w(e_m)$ , we have  $f(w) = \max(wx : x \in P_f)$ (16.18) $= \sum w(e_i) f(e_i | E_{i-1})$ (16.19) $=\sum w(e_i)(f(E_i) - f(E_{i-1}))$ (16.20)m-1 $= w(e_m)f(E_m) + \sum_{i=1}^{n} (w(e_i) - w(e_{i+1}))f(E_i)$ (16.21)
- We say that  $\emptyset \triangleq E_0 \subset E_1 \subset E_2 \subset \cdots \subset E_m = E$  forms a chain based on w.

#### • Definition of the continuous extension, once again, for reference:

 $\tilde{f}(w) = \max(wx : x \in P_f) \tag{16.22}$ 

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Min-Norm Point and SFM

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where  $\lambda_m = w(e_m)$  and otherwise  $\lambda_i = w(e_i) - w(e_{i+1})$ , where the elements are sorted according to w as before.

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where  $\lambda_m = w(e_m)$  and otherwise  $\lambda_i = w(e_i) - w(e_{i+1})$ , where the elements are sorted according to w as before.

• From convex analysis, we know  $\tilde{f}(w) = \max(wx : x \in P)$  is always convex in w for any set  $P \subseteq R^E$ , since it is the maximum of a set of linear functions (true even when f is not submodular or P is not a convex set).

#### An extension of $f_{1}$

• Recall, for any such  $w \in \mathbb{R}^E$ , we have

$$\begin{pmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \end{pmatrix} = \underbrace{(w_1 - w_2)}_{\lambda_1} \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} + \underbrace{(w_2 - w_3)}_{\lambda_2} \begin{pmatrix} 1 \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} + \underbrace{(w_{n-1} - w_n)}_{\lambda_{m-1}} \begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \\ 0 \end{pmatrix} + \underbrace{(w_m)}_{\lambda_m} \begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \\ 1 \end{pmatrix}$$
(16.25)

#### An extension of f

• Recall, for any such  $w \in \mathbb{R}^E$ , we have

$$\begin{pmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \end{pmatrix} = \underbrace{(w_1 - w_2)}_{\lambda_1} \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} + \underbrace{(w_2 - w_3)}_{\lambda_2} \begin{pmatrix} 1 \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} + \underbrace{(w_{n-1} - w_n)}_{\lambda_{m-1}} \begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \\ 0 \end{pmatrix} + \underbrace{(w_m)}_{\lambda_m} \begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \\ 1 \end{pmatrix}$$
(16.25)

• If we take w in decreasing order, then each coefficient of the vectors is non-negative (except possibly the last one,  $\lambda_m = w_m$ ).

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#### An extension of $f_1$

# • Define sets $E_i$ based on this decreasing order of w as follows, for $i=0,\ldots,n$

$$E_i \stackrel{\text{def}}{=} \{e_1, e_2, \dots, e_i\}$$
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- Note that
- Hence, from the previous and current slide, we have  $w = \sum_{i=1}^m \lambda_i \mathbf{1}_{E_i}$

#### From f back to f, even when f is not submodular

• From the continuous  $\tilde{f}$ , we can recover f(A) for any  $A \subseteq V$ .

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- From the continuous f, we can recover f(A) for any  $A \subseteq V$ .
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$$w = (w(e_1), w(e_2), \dots, w(e_m)) = (\underbrace{1, 1, 1, \dots, 1}_{|A| \text{ times}}, \underbrace{0, 0, \dots, 0}_{m-|A| \text{ times}}) \quad (16.27)$$

so that  $1_A(i) = 1$  if  $i \leq |A|$ , and  $1_A(i) = 0$  otherwise.

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• For any  $f: 2^E \to \mathbb{R}$ ,  $w = \mathbf{1}_A$ , since  $E_{|A|} = \{e_1, e_2, \dots, e_{|A|}\} = A$ :

 $\tilde{f}(w)$ 

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$$= (\mathbf{1}_{A}(|A|) - \mathbf{1}_{A}(|A|+1))f(E_{|A|}) = f(E_{|A|}) = f(A) \quad (16.29)$$

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# From $\tilde{f}$ back to f

• We can view  $\tilde{f}: [0,1]^E \to \mathbb{R}$  defined on the hypercube, with f defined as  $\tilde{f}$  evaluated on the hypercube extreme points (vertices).

Lovász extension

# From $\tilde{f}$ back to f

Min-Norm Point and SFM

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- To summarize, with  $\tilde{f}(A) = \sum_{i=1}^{m} \lambda_i f(E_i)$ , we have

$$\tilde{f}(\mathbf{1}_A) = f(A), \tag{16.30}$$

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ullet ... and when f is submodular, we also have have

$$\tilde{f}(\mathbf{1}_A) = \max\left\{\mathbf{1}_A x : x \in P_f\right\}$$
(16.31)

$$= \max \left\{ \mathbf{1}_A x : x(B) \le f(B), \forall B \subseteq E \right\}$$
(16.32)

(16.33)

Min-Norm Point and SFM

Lovász extension

#### An extension of an arbitrary $f: 2^V \to \mathbb{R}$

• Thus, for any  $f: 2^E \to \mathbb{R}$ , even non-submodular f, we can define an extension, having  $\tilde{f}(\mathbf{1}_A) = f(A), \ \forall A$ , in this way where

$$\tilde{f}(w) = \sum_{i=1}^{m} \lambda_i f(E_i)$$
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with the  $E_i = \{e_1, \ldots, e_i\}$ 's defined based on sorted descending order of w as in  $w(e_1) \ge w(e_2) \ge \cdots \ge w(e_m)$ , and where

for 
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,  $\lambda_i = \begin{cases} w(e_i) - w(e_{i+1}) & \text{if } i < m \\ w(e_m) & \text{if } i = m \end{cases}$  (16.35)

so that  $w = \sum_{i=1}^{m} \lambda_i \mathbf{1}_{E_i}$ .

Min-Norm Point and SFM

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so that  $w = \sum_{i=1}^{m} \lambda_i \mathbf{1}_{E_i}$ . •  $w = \sum_{i=1}^{m} \lambda_i \mathbf{1}_{E_i}$  is an interpolation of certain hypercube vertices.

Min-Norm Point and SFM

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so that  $w = \sum_{i=1}^{m} \lambda_i \mathbf{1}_{E_i}$ . •  $w = \sum_{i=1}^{m} \lambda_i \mathbf{1}_{E_i}$  is an interpolation of certain hypercube vertices. •  $\tilde{f}(w) = \sum_{i=1}^{m} \lambda_i f(E_i)$  is the associated interpolation of the values of f at sets corresponding to each hypercube vertex.

Min-Norm Point and SFM

#### Weighted gains vs. weighted functions

• Again sorting E descending in w, the extension summarized:

$$\tilde{f}(w) = \sum_{i=1}^{m} w(e_i) f(e_i | E_{i-1})$$

$$= \sum_{i=1}^{m} w(e_i) (f(E_i) - f(E_{i-1}))$$

$$= w(e_m) f(E_m) + \sum_{i=1}^{m-1} (w(e_i) - w(e_{i+1})) f(E_i)$$

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Min-Norm Point and SFM

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• Again sorting E descending in w, the extension summarized:

$$= w(e_m)f(E_m) + \sum_{i=1}^{m-1} (w(e_i) - w(e_{i+1}))f(E_i)$$
 (16.38)

$$=\sum_{i=1}^{m}\lambda_{i}f(E_{i})$$
(16.39)

• So  $\tilde{f}(w)$  seen either as sum of weighted gain evaluatiosn (Eqn. (16.36), or as sum of weighted function evaluations (Eqn. (16.39)).

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# The Lovász extension of $f: 2^E \to \mathbb{R}$

• Lovász showed that if a function  $\tilde{f}(w)$  defined as in Eqn. (16.34) is convex, then f must be submodular.

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- Note, also possible to define this when  $f(\emptyset) \neq 0$  (but doesn't really add any generality).

#### Theorem 16.4.1

A function  $f: 2^E \to \mathbb{R}$  is submodular iff its Lovász extension  $\tilde{f}$  of f is convex.

#### Proof.

• We've already seen that if f is submodular, its extension can be written via Eqn.(16.34) due to the greedy algorithm, and therefore is also equivalent to  $\tilde{f}(w) = \max \{wx : x \in P_f\}$ , and thus is convex.

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- Conversely, suppose the Lovász extension  $\tilde{f}(w) = \sum_i \lambda_i f(E_i)$  of some function  $f: 2^E \to \mathbb{R}$  is a convex function.
- We note that, based on the extension definition, in particular the definition of the  $\{\lambda_i\}_i$ , we have that  $\tilde{f}(\alpha w) = \alpha \tilde{f}(w)$  for any  $\alpha \in \mathbb{R}_+$ . I.e., f is a positively homogeneous convex function.

. .

Min-Norm Point and SFM

### ... proof of Thm. 16.4.1 cont.

• Earlier, we saw that  $\tilde{f}(\mathbf{1}_A) = f(A)$  for all  $A \subseteq E$ .

Min-Norm Point and SFM

- Earlier, we saw that  $\tilde{f}(\mathbf{1}_A) = f(A)$  for all  $A \subseteq E$ .
- Now, given  $A, B \subseteq E$ , we will show that  $\tilde{f}(\mathbf{1}_A + \mathbf{1}_B) = \tilde{f}(\mathbf{1}_{A \cup B} + \mathbf{1}_{A \cap B})$  (16.40)  $= f(A \cup B) + f(A \cap B).$  (16.41)

# Lovász Extension, Submodularity and Convexity

- Earlier, we saw that  $f(\mathbf{1}_A) = f(A)$  for all  $A \subseteq E$ .
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- Let  $C = A \cap B$ , order E based on decreasing  $w = \mathbf{1}_A + \mathbf{1}_B$  so that  $w = (w(e_1), w(e_2), \dots, w(e_m))$  (16.42)  $= (\underbrace{2, 2, \dots, 2}_{i \in C}, \underbrace{1, 1, \dots, 1}_{i \in A \triangle B}, \underbrace{0, 0, \dots, 0}_{i \in E \setminus (A \cup B)})$  (16.43)

## Lovász Extension, Submodularity and Convexity

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- Then, considering  $\tilde{f}(w) = \sum_i \lambda_i f(E_i)$ , we have  $\lambda_{|C|} = 1$ ,  $\lambda_{|A \cup B|} = 1$ , and  $\lambda_i = 0$  for  $i \notin \{|C|, |A \cup B|\}$ .

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- Let  $C = A \cap B$ , order E based on decreasing  $w = \mathbf{1}_A + \mathbf{1}_B$  so that  $w = (w(e_1), w(e_2), \dots, w(e_m))$  (16.42)  $= (\underbrace{2, 2, \dots, 2}_{i \in C}, \underbrace{1, 1, \dots, 1}_{i \in A \triangle B}, \underbrace{0, 0, \dots, 0}_{i \in E \setminus (A \cup B)})$  (16.43)
- Then, considering  $\tilde{f}(w) = \sum_{i} \lambda_i f(E_i)$ , we have  $\lambda_{|C|} = 1$ ,  $\lambda_{|A \cup B|} = 1$ , and  $\lambda_i = 0$  for  $i \notin \{|C|, |A \cup B|\}$ .
- But then  $E_{|C|} = A \cap B$  and  $E_{|A \cup B|} = A \cup B$ . Therefore,  $\tilde{f}(w) = \tilde{f}(\mathbf{1}_A + \mathbf{1}_B) = f(A \cap B) + f(A \cup B)$ .

Min-Norm Point and SFM

#### ... proof of Thm. 16.4<u>.1 cont.</u>

• Also, since  $\tilde{f}$  is convex (by assumption) and positively homogeneous, we have for any  $A, B \subseteq E$ ,

 $0.5[f(A\cap B)+f(A\cup B)]$ 

Min-Norm Point and SFM

### ... proof of Thm. 16.4.1 cont.

• Also, since  $\tilde{f}$  is convex (by assumption) and positively homogeneous, we have for any  $A, B \subseteq E$ ,

$$0.5[f(A \cap B) + f(A \cup B)] = 0.5[\tilde{f}(\mathbf{1}_A + \mathbf{1}_B)]$$
(16.44)

Min-Norm Point and SFM

### ... proof of Thm. 16.4.1 cont.

• Also, since  $\tilde{f}$  is convex (by assumption) and positively homogeneous, we have for any  $A, B \subseteq E$ ,

$$0.5[f(A \cap B) + f(A \cup B)] = 0.5[\tilde{f}(\mathbf{1}_A + \mathbf{1}_B)]$$
(16.44)

 $= f(0.5\mathbf{1}_A + 0.5\mathbf{1}_B) \tag{16.45}$ 

Min-Norm Point and SFM

### ... proof of Thm. 16.4.1 cont.

• Also, since  $\tilde{f}$  is convex (by assumption) and positively homogeneous, we have for any  $A, B \subseteq E$ ,

$$0.5[f(A \cap B) + f(A \cup B)] = 0.5[\tilde{f}(\mathbf{1}_A + \mathbf{1}_B)]$$
(16.44)

$$=\tilde{f}(0.5\mathbf{1}_A+0.5\mathbf{1}_B)$$
 (16.45)

$$\leq 0.5 \tilde{f}(\mathbf{1}_A) + 0.5 \tilde{f}(\mathbf{1}_B)$$
 (16.46)

Min-Norm Point and SFM

#### ... proof of Thm. 16.4.1 cont.

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$$\leq 0.5\tilde{f}(\mathbf{1}_A) + 0.5\tilde{f}(\mathbf{1}_B)$$
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$$= 0.5(f(A) + f(B))$$
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Min-Norm Point and SFM

#### ... proof of Thm. 16.4.1 cont.

• Also, since  $\tilde{f}$  is convex (by assumption) and positively homogeneous, we have for any  $A, B \subseteq E$ ,

$$0.5[f(A \cap B) + f(A \cup B)] = 0.5[\tilde{f}(\mathbf{1}_A + \mathbf{1}_B)]$$
(16.44)

$$=\tilde{f}(0.5\mathbf{1}_A+0.5\mathbf{1}_B)$$
 (16.45)

$$\leq 0.5\tilde{f}(\mathbf{1}_{A}) + 0.5\tilde{f}(\mathbf{1}_{B}) \qquad (16.46)$$

$$= 0.5(f(A) + f(B))$$
 (16.47)

• Thus, we have shown that for any  $A, B \subseteq E$ ,

 $f(A \cup B) + f(A \cap B) \le f(A) + f(B)$  (16.48)

so f must be submodular.

EE596b/Spring 2014/Submodularity - Lecture 16 - May 21st, 2014

### Edmonds - Submodularity - 1969

#### SUBMODULAR FUNCTIONS, MATROIDS, AND CERTAIN POLYHEDRA\*

Jack Edmonds

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#### I.

The viewpoint of the subject of matroids, and related areas of lattice theory, has always been, in one way or another, abstraction of algebraic dependence or, equivalently, abstraction of the incidence relations in geometric representations of algebra. Often one of the

### Lovász - Submodularity - 1983

### Submodular functions and convexity

#### L. Lovász

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#### 0. Introduction

In "continuous" optimization convex functions play a central role. Besides elementary tools like differentiation, various methods for finding the minimum of a convex function constitute the main body of nonlinear optimization. But even linear programming may be viewed as the optimization of very special (linear) objective functions over very special convex domains (polyhedra). There are several reasons for this popularity of convex functions:

- Convex functions occur in many mathematical models in economy, engineering, and other sciencies. Convexity is a very natural property of various functions and domains occuring in such models; quite often the only non-trivial property which can be stated in general.