## EE512A – Advanced Inference in Graphical Models

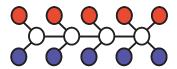
— Fall Quarter, Lecture 18 —

http://j.ee.washington.edu/~bilmes/classes/ee512a\_fall\_2014/

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- Wainwright and Jordan Graphical Models, Exponential Families, and Variational Inference http://www.nowpublishers.com/product. aspx?product=MAL&doi=2200000001
- Should have read chapters 1 through 5 in our book. Read chapter 7.
- Also read chapter 8 (integer programming, although we probably won't cover that chapter in class unfortunately).
- Also should have read "Divergence measures and message passing" by Thomas Minka, and "Structured Region Graphs: Morphing EP into GBP", by Welling, Minka, and Teh.
- Assignment due Wednesday (Dec 3rd) night, 11:45pm. Final project proposal final progress report (one page max).
- Update: For status update, final writeup, and talk, use notation as close as possible to that used in class!

- L1 (9/29): Introduction, Families, Semantics
- L2 (10/1): MRFs, elimination, Inference on Trees
- L3 (10/6): Tree inference, message passing, more general queries, non-tree)
- L4 (10/8): Non-trees, perfect elimination, triangulated graphs
- L5 (10/13): triangulated graphs, k-trees, the triangulation process/heuristics
- L6 (10/15): multiple queries, decomposable models, junction trees
- L7 (10/20): junction trees, begin intersection graphs
- L8 (10/22): intersection graphs, inference on junction trees
- L9 (10/27): inference on junction trees, semirings,
- L10 (11/3): conditioning, hardness, LBP

- L11 (11/5): LBP, exponential models,
- L12 (11/10): exponential models, mean params and polytopes,
- L13 (11/12): polytopes, tree outer bound, Bethe entropy approx.
- L14 (11/17): Bethe entropy approx, loop series correction
- L15 (11/19): Hypergraphs, posets, Mobius, Kikuchi
- L16 (11/24): Kikuchi, Expectation Propagation
- L17 (11/26): Expectation Propagation, Mean Field
- L18 (12/1): Structured mean field, Convex relaxations and upper bounds, tree reweighted case
- L19 (12/3):
- Final Presentations: (12/10):

# Conjugate Duality, Maximum Likelihood, Negative Entropy

## Theorem 18.2.3 (Relationship between A and $A^*$ )

(a) For any  $\mu \in \mathcal{M}^{\circ}$ ,  $\theta(\mu)$  unique canonical parameter sat. matching condition, then conj. dual takes form:

$$A^{*}(\mu) = \sup_{\theta \in \Omega} \left( \langle \theta, \mu \rangle - A(\theta) \right) = \begin{cases} -H(p_{\theta(\mu)}) & \text{if } \mu \in \mathcal{M}^{\circ} \\ +\infty & \text{if } \mu \notin \overline{\mathcal{M}} \end{cases}$$
(18.3)

(b) Partition function has variational representation (dual of dual)

$$A(\theta) = \sup_{\mu \in \mathcal{M}} \left\{ \langle \theta, \mu \rangle - A^*(\mu) \right\}$$
 (18.4)

(c) For  $\theta \in \Omega$ , sup occurs at  $\mu \in \mathcal{M}^{\circ}$  of moment matching conditions

$$\mu = \int_{D_X} \phi(x) p_{\theta}(x) \nu(dx) = \mathbb{E}_{\theta}[\phi(X)] = \nabla A(\theta)$$
 (18.5)

## Variational Approach Amenable to Approximation

Original variational representation of log partition function

$$A(\theta) = \sup_{\mu \in \mathcal{M}} \{ \langle \theta, \mu \rangle - A^*(\mu) \}$$
 (18.1)

where dual takes form:

$$A^{*}(\mu) = \sup_{\theta \in \Omega} (\langle \theta, \mu \rangle - A(\theta)) = \begin{cases} -H(p_{\theta(\mu)}) & \text{if } \mu \in \mathcal{M}^{\circ} \\ +\infty & \text{if } \mu \notin \overline{\mathcal{M}} \end{cases}$$
 (18.2)

- Given efficient expression for  $A(\theta)$ , we can compute marginals of interest.
- Above expression (dual of the dual) offers strategies to approximate or (upper or lower) bound  $A(\theta)$ . We either approximate  $\mathcal{M}$  or  $-A^*(\mu)$  or (most likely) both.

## Variational Approximations we cover

- Set  $\mathcal{M} \leftarrow \mathbb{L}$  and  $-A^*(\mu) \leftarrow H_{\mathsf{Bethe}}(\tau)$  to get Bethe variational approximation, LBP fixed point.
- ② Set  $\mathcal{M} \leftarrow \mathbb{L}_t(G)$  (hypergraph marginal polytope),  $-A^*(\mu) \leftarrow H_{\mathsf{app}}(\tau)$  where  $H_{\mathsf{app}} = \sum_{g \in E} c(g) H_g(\tau_g)$  (via Möbius) to get Kikuchi variational approximation, message passing on hypergraphs.
- **9** Partition  $\tau$  into  $(\tau, \tilde{\tau})$ , and set  $\mathcal{M} \leftarrow \mathcal{L}(\phi, \Phi)$  and set  $-A^*(\mu) \leftarrow H_{ep}(\tau, \tilde{\tau})$  to get expectation propagation.
- **Mean field** (from variational perspective) is (with  $\mathcal{M}_F(G) \subseteq \mathcal{M}$ ) **l.b.**:

$$A(\theta) \ge \max_{\mu \in \mathcal{M}_F(G)} \{ \langle \mu, \theta \rangle - A_F^*(\mu) \} = A_{\mathsf{mf}}(\theta) \tag{18.1}$$

## EP as variational: Summary of key points

- Fixed points of EP exist assuming Lagrangian form has at least one optimum.
- No guarantees that EP will converge, but if it does it will be at a stationary point of the Lagrangian.
- EP can be seen to be based on variational framework, using Bethe-like entropy and convex outer bound for the mean parameters.
- When base distribution is unaries and  $\Phi^i$  is the edges of a graph, we in fact get standard Bethe approximation, and standard sum-product LBP.
- Moment matching of EP can be seen as striving for solution of associated Lagrangian.
- Lost of flexibility here, depending on what the base distribution is (e.g., could be a k-tree, clusters, or many other structures as well).
- Can also be done for Gaussian mixture and other distributions.
  Many more details, variations, and possible roads to new research. See text and also see Tom Minka's papers. http://research.microsoft.com/en-us/um/people/minka/papers/

### Mean Field

- ullet So far, we have been using an outer bound on  ${\mathcal M}.$
- In mean-field methods, we use an "inner bound", a subset of  $\mathcal M$  constructed so as to make the optimization of  $A(\theta)$  easier.
- Since subset, we get immediate bound on  $A(\theta)$ , all else (i.e., the entropy) being equal.
- Key: we based the inner bound on a "tractable family" like a 1-tree or even a 0-tree (all independent) so that the variational problem can be computed efficiently.
- Convexity of the optimization problem is often lost still, however, in the general case (due to the inner bound).
- Thus, in mean field, we will get a lower bound on  $A(\theta)$  but not a convex procedure to find it (both good and bad news).

## Tractable Families (for mean field approach)

- We have graph G = (V, E) which is intractable and we find a spanning subgraph (recall, spanning = all nodes, subgraph = subset of edges), i..e,  $F = (V, E_F)$  where  $E_F \subseteq E$ .
- Simplest example:  $F = (V, \emptyset)$  all independence model.
- Tree example:  $F = (V, E_T)$  where edges  $E_T \subset E$  constitute a spanning tree.
- Exponential family, sufficient statistics  $\phi = (\phi_{\alpha}, \alpha \in \mathcal{I})$  associated with this family  $\mathcal{I}(F) \subseteq \mathcal{I}$ . These are the statistics that need respect the Markov properties of only the subgraph F.
- ullet  $\Omega$  gets smaller too, canonical F-respecting parameters are of the form:

$$\mathbb{R}^{|\mathcal{I}|} \ni \Omega(F) \triangleq \{ \theta \in \Omega | \theta_{\alpha} = 0 \ \forall \alpha \in \mathcal{I} \setminus \mathcal{I}(F) \} \subseteq \Omega.$$
 (18.14)

Notice, all parameters associated with sufficient statistic not in  $\mathcal{I}(F)$  are set to zero, those statistics are nonexistent in F.

ullet If parameter was not zero, model would not respect the familiy of F.

## Inner bound Approximate Polytope

- Before, we had  $\mathcal{M}(G;\phi)(=\mathcal{M}_G(G;\phi))$ , all possible mean parameters associated with G and associated set of sufficient statistics  $\phi$ .
- ullet For a given subgraph F, we only consider those mean parameters possible under F-respecting models. I.e.,

$$\mathcal{M}_F(G;\phi) = \left\{ \mu \in \mathbb{R}^d | \mu = \mathbb{E}_{\theta}[\phi(x)] \text{ for some } \theta \in \Omega(F) \right\}$$
 (18.18)

• Therefore, since  $\theta \in \Omega(F) \subseteq \Omega$ , we have that

$$\mathcal{M}_F^{\circ}(G;\phi) \subseteq \mathcal{M}^{\circ}(G;\phi) \tag{18.19}$$

and so  $\mathcal{M}_F^\circ(G;\phi)$  is an inner approximation of the set of realizable mean parameters.

• Shorthand notation:  $M_F^{\circ}(G) = M_F^{\circ}(G;\phi)$  and  $M^{\circ}(G) = M^{\circ}(G;\phi)$ 

### Tractable Dual

- Normally dual  $A^*(\mu) = \sup_{\theta \in \Omega} (\langle \theta, \mu \rangle A(\theta))$  is intractable or unavailable, but key idea is that if  $\mu \in \mathcal{M}_F(G)$  it will be possible to compute easily.
- Thus, goal of mean field (from variational approximation perspective) is to form  $A_{\rm MF}(\theta)$  where:

$$A(\theta) \ge \max_{\mu \in \mathcal{M}_F(G)} \left\{ \langle \mu, \theta \rangle - A_F^*(\mu) \right\} \triangleq A_{\mathsf{MF}}(\theta) \tag{18.23}$$

where  $A_F^*(\mu)$  corresponds to dual function restricted to inner bound set  $\mathcal{F}(G)$ . I.e., when we expand  $A_F^*(\mu)$ , we can take advantage of the fact that  $\mu$  is restricted in all cases, so  $A_F^*(\mu)$  might be greatly simplified relative to  $A^*(\mu)$ .

• Note, for  $\mu \in \mathcal{M}_F(G)$  and since  $\mathcal{M}_F(G) \subseteq \mathcal{M}(G)$ ,  $A_F^*(\mu)$  is not an approximation, rather it is just easy to compute.

## Mean field, KL-Divergence, Exponential Model Families

• Thus, solving the mean-field variational problem (see Eqn. (??)) of:

$$\max_{\mu \in \mathcal{M}_F(G)} \{ \langle \mu, \theta \rangle - A_F^*(\mu) \} = \max_{\mu \in \mathcal{M}_F(G)} \{ \langle \mu, \theta \rangle - A^*(\mu) \}$$
 (18.34)

is identical to minimizing KL Divergence  $D(\mu||\theta)$  subject to constraint  $\mu \in \mathcal{M}_F(G)$ .

• I.e., mean field can be seen as finding the best approximation, in terms of this particular KL-divergence, to  $p_{\theta}$ , over a family of "nice" distributions  $M_F(G)$ .

## Naïve Mean field for Ising Model: optimization

• We get variational lower bound problem

$$A(\theta) \ge \max_{(\mu_1, \dots, \mu_m) \in [0,1]^m} \left\{ \sum_{s \in V} \theta_s \mu_s + \sum_{(s,t) \in E} \theta_{st} \mu_s \mu_t + \sum_{s \in V} H_s(\mu_s) \right\}$$
(18.35)

- ullet Have constrained form of edge mean parameters  $\mu_{st}=\mu_s\mu_t$
- $(\mu_1, \ldots, \mu_m) \in [0,1]^m$  is m-D hypercube.
- We have a non-convex problem, so while it is a bound, it might be hard to get as tight as possible.
- One way to optimize is to do coordinate ascent (given otherwise fixed vector, optimize one value at a time).
- If each coordinate optimization is optimal, we'll get a stationary point.
- Fortunately, each coordinate optimization is concave!

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- In Structured mean field, we exploit this and it again can be seen in our variational framework
- We first see a nice way that we can use fixed points of the mean field primal/dual equations to derive a general form of the mean field update.

• Again,  $\mathcal{I}(F)$  is set of suff. stats. corresponding to F, and we have corresponding mean vector  $\mu(F) = (\mu_{\alpha}, \alpha \in \mathcal{I}(F))$ .

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- Key thing: in mean field,  $\mu(F) \in \mathcal{M}(F)$  and there is no real need to mention the full  $M_F(G)$ . Also, the dual  $A_F^*$  depends on only  $\mu(F)$  not  $\mu$  (the other values are derivations from entries within  $\mu(F)$ .

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- Other mean parameters  $\mu_{\beta}$  for  $\beta \in \mathcal{I} \setminus \mathcal{I}(F)$  do play a role in the value of the mean field variational problem but their value is derivable from values  $\mu(F)$ , thus we can express the  $\mu_{\beta}$  in functional form based on values  $\mu(F)$ .

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- Ex: mean field Ising, edges  $(s,t) \in E$ , get  $\mu_{st} = g_{st}(\mu(F)) = \mu_s \mu_t$ .

• The mean field optimization problem becomes

$$\max_{\mu \in \mathcal{M}_{F}(G)} \left\{ \langle \mu, \theta \rangle - A_{F}^{*}(\mu) \right\}$$

$$= \max_{\mu(F) \in \mathcal{M}(F)} \left\{ \underbrace{\sum_{\alpha \in \mathcal{I}(F)} \theta_{\alpha} \mu_{\alpha} + \sum_{\alpha \in \mathcal{I}^{c}(F)} \theta_{\alpha} g_{\alpha}(\mu(F)) - A_{F}^{*}(\mu(F))}_{f(\mu(F))} \right\}$$

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• With this, we can recover our sigmoid mean field coordinate update process by iterating fixed point equations of f, i.e., for  $\beta \in \mathcal{I}(F)$ ,

$$\frac{\partial f}{\partial \mu_{\beta}}(\mu(F)) = \theta_{\beta} + \sum_{\alpha \in \mathcal{I}(G) \setminus \mathcal{I}(F)} \theta_{\alpha} \frac{\partial g_{\alpha}}{\partial \mu_{\beta}}(\mu(F)) - \frac{\partial A_{F}^{*}}{\partial \mu_{\beta}}(\mu(F))$$
 (18.3)

• Setting this to zero, and then aggregating/concatenating over  $\beta \in \mathcal{I}(F)$ , vector fix point condition is:

$$\nabla A_F^*(\mu(F)) = \theta + \sum_{\alpha \in \mathcal{I}(G) \setminus \mathcal{I}(F)} \theta_\alpha \nabla g_\alpha(\mu(F))$$
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•  $\nabla A$  is the forward mapping, maps from canonical to mean parameters, and  $\nabla A^*$  does the reverse. Hence, naming  $\gamma(F) = \nabla A(\mu(F))$ , gives a parameter update equation for  $\beta \in \mathcal{I}(F)$ 

$$\gamma_{\beta}(F) \leftarrow \theta_{\beta} + \sum_{\alpha \in \mathcal{I}(G) \setminus \mathcal{I}(F)} \theta_{\alpha} \frac{\partial g_{\alpha}}{\partial \mu_{\beta}}(\mu(F))$$
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• Above is the mean field update, mapping from canonical parameters  $(\theta_{\beta}, \text{ and } \theta_{\alpha} \text{ for } \alpha \in \mathcal{I}(G) \setminus \mathcal{I}(F))$  and using the mean parameters  $\mu(F)$  to new updated canonical parameters  $\gamma_{\beta}(F)$  for  $\beta \in \mathcal{I}(F)$ . It is to be repeated over and over.

Str. Mean Field

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- Since we're using a tractable sub-structure F, we can then update the out-of-date mean parameters using any exact inference algorithm (e.g., junction tree, possible since sub-structure is tractable), and then repeat Eqn. (18.5).

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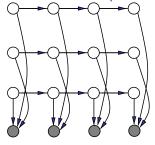
• This generalizes our mean field coordinate ascent update from before, where in that case we would get  $\frac{\partial A_F}{\partial \gamma_s}$  as being the sigmoid mapping.

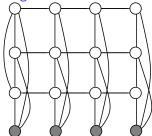
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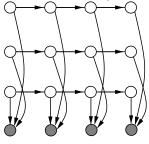
- This generalizes our mean field coordinate ascent update from before, where in that case we would get  $\frac{\partial A_F}{\partial \gamma_g}$  as being the sigmoid mapping.
- But here, we can use this for any tractable substructure (e.g., trees or chains or collections thereof).

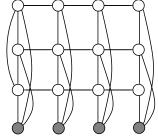
• This idea was developed and applied using factorial HMMs.





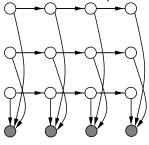
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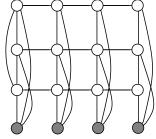




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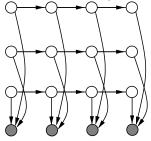
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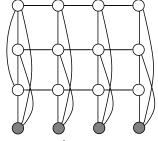




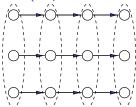
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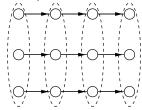




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- After moralization, covering hypergraph consists of tractable sub-substructure hyperedges  $F = \left\{ \left\{ x_t^i, x_{t+1}^i \right\} : i \in [M], t \in [T] \right\}$  and remaining structure  $E \setminus F = \left\{ \left\{ x_t^1, x_t^2, \dots, x_t^M \right\} : t \in [T] \right\}$ .



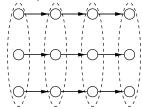
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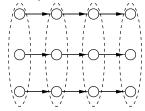
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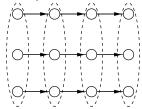


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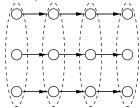
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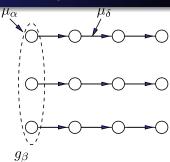
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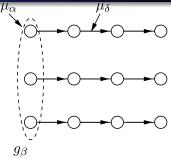
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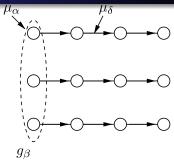
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• Under this independent chains case, we have that for each  $\beta \in \mathcal{I} \setminus \mathcal{I}(F)$ , derivable functions have form  $g_{\beta}(\mu(F)) = \prod_{i=1}^M f_i(\{\mu_i(F)\})$ , for some functions  $f_i$ . This is fully factored, so is easy to work with, maintains separate chains.

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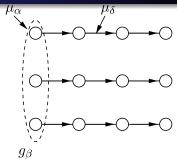
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- Each update of form Eqn. (18.5) updates parameters for  $\beta \in \mathcal{I}(F)$ , corresponds to all edges of all M Markov chains.
- To recover mean parameters (or do Eqn. (18.6)), need only forward-backward procedure on each chain separately,  $O(MTr^2)$ .

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- ullet In this next chapter (Chap 7), we will "convexify"  $H(\mu)$  and at the same time produce upper bounds.

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- As before,  $\mathcal{M}(F)$  is set of realizable mean parameters associated with F, and  $\mu(F) \in \mathcal{M}(F)$ . Thus,  $\mathcal{M}(F) \subseteq \mathbb{R}^{|\mathcal{I}(F)|}$ , and

$$\mathcal{M}(F) = \left\{ \mu \in \mathbb{R}^{|\mathcal{I}(F)|} | \exists p \text{ s.t. } \mu_{\alpha} = \mathbb{E}_{p}[\phi_{\alpha}(X)] \ \forall \alpha \in \mathcal{I}(F) \right\}$$
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Note  $\mathcal{M}_F(G) \neq \mathcal{M}(F)$ .

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- Critically, we have that  $H(\mu(F)) \geq H(\mu) = H(p_{\mu})$ , as we show next.

#### Proposition 18.4.1 (Maximum Entropy Bounds)

Given any mean parameter  $\mu \in \mathcal{M}$  and its projection  $\mu(F)$  onto any subgraph F, we have the bound

$$A^*(\mu(F)) \le A^*(\mu) \tag{18.9}$$

or alternatively stated,  $H(\mu(F)) \ge H(\mu)$ , entropy of projection is higher.

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- Thus,  $H(\mu(F)) > H(\mu)$ .

#### Proof.

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Tree Re-weighted Case

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- ullet This so far gives us an upper bound on  $A(\theta)$ , but we still need an outer bound. The combination will give us our uppper bound on  $A(\theta)$ .

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$$\mathcal{L}(G; \mathfrak{D}) = \left\{ \tau \in \mathbb{R}^d | \tau(F) \in \mathcal{M}(F) \ \forall F \in \mathfrak{D} \right\}$$
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- ullet Also note,  $\mathcal{L}(G;\mathfrak{D})$  is convex since it is the intersection of a set of convex sets.

# Convex Upper Bounds

ullet Combining the upper bound on entropy, and the outer bound on  $\mathcal{M}$ , we get a new variational approximation to the cumulant function.

$$B_{\mathfrak{D}}(\theta; \rho) \stackrel{\Delta}{=} \sup_{\tau \in \mathcal{L}(G; \mathfrak{D})} \left\{ \langle \tau, \theta \rangle + \sum_{F \in \mathfrak{D}} \rho(F) H(\tau(F)) \right\}$$
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# Convex Upper Bounds

• Combining the upper bound on entropy, and the outer bound on  $\mathcal{M}$ , we get a new variational approximation to the cumulant function.

$$B_{\mathfrak{D}}(\theta; \rho) \stackrel{\Delta}{=} \sup_{\tau \in \mathcal{L}(G; \mathfrak{D})} \left\{ \langle \tau, \theta \rangle + \sum_{F \in \mathfrak{D}} \rho(F) H(\tau(F)) \right\}$$
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- Evaluating the objective (optimization) is concave, so possible to get!
- Also,  $\mathcal{L}(G;\mathfrak{D})$  is a convex outer bound on  $\mathcal{M}(G)$
- Thus  $B_{\mathfrak{D}}(\theta; \rho)$  is convex, has a global optimal solution, it approximates  $A(\theta)$ , and best of all is an upper bound,  $A(\theta) \leq B_{\mathfrak{D}}(\theta; \rho)$

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ullet We want to use this to see what happens when we take the expected value w.r.t. distribution ho.

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- The expression becomes

$$H(\mu) \le \sum_{s \in V} H_s(\mu_s) - \sum_{(s,t) \in E} \rho_{st} I_{st}(\mu_{st})$$
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Note right hand sum is over all E (not just a given spanning tree) and terms are weighted by probability of the given edge  $\rho_{st}$ .

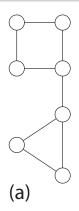
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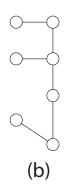
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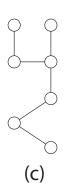
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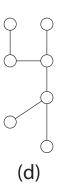
•  $\rho_{st}$  is edge appearance probability,  $\rho=(\rho_{st},(s,t)\in E)$  is spanning tree polytope.

## Edge appearance probabilities example



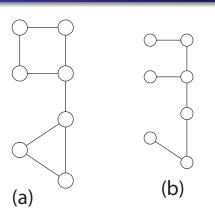


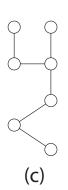


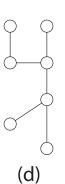


• (a) a graph 
$$G = (V, E)$$
 with  $m = |V| = 7$ 

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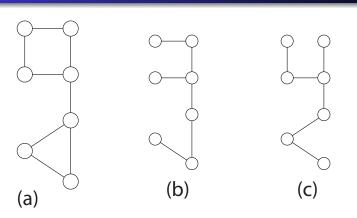






- (a) a graph G = (V, E) with m = |V| = 7
- (b), (c), and (d) various spanning trees, each with probability 1/3.

## Edge appearance probabilities example



- (a) a graph G=(V,E) with m=|V|=7
- $\bullet$  (b), (c), and (d) various spanning trees, each with probability 1/3.
- ullet What are the edge appearance probabilities  $ho_{st}$ ?

 $\bullet$  We also need outer bound on  $\mathcal{M}$ .

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- In Bethe case, however, we did not have a bound on entropy, only an outer bound on the marginal polytope. Now, however, we also have a (convexification based) bound on entropy.

### Theorem 18.5.1 (Tree-Reweighted Bethe and Sum-Product)

(a) For any choice of edge appearance vector  $\rho = (\rho_{st}, (s,t) \in E)$  in the spanning tree polytope, the cumulant function  $A(\theta)$  evaluated at  $\theta$  is upper bounded by the solution of the tree reweighted Bethe variational problem (BVP):

$$B_{\mathfrak{T}}(\theta; \rho) = \max_{\tau \in \mathbb{L}(G)} \left\{ \langle \tau, \theta \rangle + \sum_{s \in V} H_s(\tau_s) - \sum_{(s,t) \in E} \rho_{st} I_{st}(\tau_{st}) \right\}$$

$$\geq A(\theta)$$
(18.19)

For any edge appearance vector such that  $\rho_{st} > 0$  for all edges (s,t), this problem is strictly convex with a unique optimum.

#### Theorem 18.5.1 (Tree-Reweighted Bethe and Sum-Product)

(b) The tree-reweighted BVP can be solved using the tree-reweighted sum-product updates

$$M_{t \to s}(x_s) \leftarrow \kappa \sum_{x_t' \in \mathcal{X}_t} \varphi_{st}(x_s, x_t') \frac{\prod_{v \in N(t) \setminus \{s\}} \left[ M_{v \to t}(x_t') \right]^{\rho_{vt}}}{\left[ M_{s \to t}(x_t') \right]^{(1 - \rho_{ts})}} \quad (18.21)$$

where  $\varphi_{st}(x_s, x_t') = \exp\left(\frac{1}{\rho_{st}}\phi_{st}(x_s, x_t') + \theta_t(x_t')\right)$ . The updates have a unique fixed point under assumptions given in (a).

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$$\rho \in \operatorname{conv}(\{\mathbf{1}_T : T \in \mathfrak{T}\}) \tag{18.22}$$

where  $conv(\cdot)$  is the convex hull of its argument.

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$$P_r = \left\{ x \in \mathbb{R}_+^E : x(A) \le r(A), \forall A \subseteq E \right\}$$
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- Then if T is a spanning tree,  $\mathbf{1}_T \in B_r$ , and  $B_r = \operatorname{conv}(\{\mathbf{1}_T : T \in \mathfrak{T}\})$ .
- Edmonds showed that a simple fast greedy procedure will maximize a linear function over this polytope, and this can be useful for finding good points in the spanning tree polytope.

## Tree-reweighted sum-product: convex vs. upper bound

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#### Tree-reweighted sum-product: convex vs. upper bound

- In above case, we have both a convexification of the cumulant and an upper bound property.
- It should be pointed out that these are not mutual requirements: one can have convex without upper bound and vice verse.

#### Tree-reweighted sum-product fixed point

The fixed point we ultimately reach has following form:

$$\tau_s^*(x_s) = \kappa \exp\{\theta_s(x_s)\} \prod_{v \in N(s)} [M_{v \to s}^*(x_s)]^{\rho_{vs}}$$
 (18.26)

$$\tau_{st}^*(x_s, x_t) = \kappa \varphi_{st}(x_s, x_t) \frac{\prod_{v \in N(s) \setminus t} [M_{vs}^*(x_s)]^{\rho_{vs}} \prod_{v \in N(t) \setminus s} [M_{vt}^*(x_t)]^{\rho_{vt}}}{[M_{ts}^*(x_s)]^{(1-\rho_{st})} [M_{st}^*(x_t)]^{(1-\rho_{ts})}}$$
(18.27)

with  $\varphi_{st}(x_s, x_t) = \exp\left\{\frac{1}{\rho_{st}}\theta_{st}(x_s, x_t) + \theta_s(x_s) + \theta_t(x_t)\right\}$  where the \* versions are the final (convergent) messages.

 $\bullet$  In practice: damping of messages M appears in practice to help reach convergence, where each new message is a convex mixture of the previous version of itself and the new message according to the equations.

Str. Mean Field

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This leads to a convexified Kikuchi variational problem

$$A(\theta) \le B_{\mathfrak{B}(t)}(\theta; \rho) = \max_{\tau \in \mathbb{L}(G)} \left\{ \langle \tau, \theta \rangle + \mathbb{E}_{\rho}[H(\tau(T))] \right\}$$
 (18.29)

same form (but different than) before.

• Optimizing  $\rho$  over hypertree polytope is hard, unfortunately.

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- Convexified forms of EP

$$H_{\rm ep}(\tau, \tilde{\tau}; \rho) = H(\tau) + \sum_{\ell=1}^{d_I} \rho(\ell) [H(\tau, \tilde{\tau}^{\ell}) - H(\tau)]$$
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where 
$$\sum_{\ell} \rho(\ell) = 1$$
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- In this case, reweighted entropy is concave!
- Lagrangian formulation leads to solutions that are a form of "reweighted" EP, ideas which also are sometimes called "power EP" (blending the above reweighted sum-product ideas and EP).

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  - use of conditional entropy, various forms of use of polyhedral approximations.
- This is still an active research area!

# Variational Approximations we cover

- Set  $\mathcal{M} \leftarrow \mathbb{L}$  and  $-A^*(\mu) \leftarrow H_{\mathsf{Bethe}}(\tau)$  to get Bethe variational approximation, LBP fixed point.
- ② Set  $\mathcal{M} \leftarrow \mathbb{L}_t(G)$  (hypergraph marginal polytope),  $-A^*(\mu) \leftarrow H_{\mathsf{app}}(\tau)$  where  $H_{\mathsf{app}} = \sum_{g \in E} c(g) H_g(\tau_g)$  (via Möbius) to get Kikuchi variational approximation, message passing on hypergraphs.
- **③** Partition  $\tau$  into  $(\tau, \tilde{\tau})$ , and set  $\mathcal{M} \leftarrow \mathcal{L}(\phi, \Phi)$  and set  $-A^*(\mu) \leftarrow H_{\text{ep}}(\tau, \tilde{\tau})$  to get expectation propagation.
- **Mean field** (from variational perspective) is (with  $\mathcal{M}_F(G) \subseteq \mathcal{M}$ ) **l.b.**:

$$A(\theta) \ge \max_{\mu \in \mathcal{M}_F(G)} \left\{ \langle \mu, \theta \rangle - A_F^*(\mu) \right\} = A_{\mathsf{mf}}(\theta) \tag{18.1}$$

① Upper bound Convexified/tree reweighted LBP, entropy upper bounds  $H(\tau(F))$  for all members  $F \in \mathfrak{D}$  of tractable substructures. Get **U.b.**:

$$A(\theta) \le B_{\mathfrak{D}}(\theta; \rho) \stackrel{\Delta}{=} \sup_{\tau \in \mathcal{L}(G; \mathfrak{D})} \left\{ \langle \tau, \theta \rangle + \sum_{F \in \mathfrak{D}} \rho(F) H(\tau(F)) \right\}$$
 (18.2)

with  $\mathcal{L}(G;\mathfrak{D}) = \bigcap_{F \in \mathfrak{D}} \mathcal{M}(F)$ 

# Sources for Today's Lecture

 Wainwright and Jordan Graphical Models, Exponential Families, and Variational Inference http://www.nowpublishers.com/product. aspx?product=MAL&doi=220000001

Str. Mean Field